



Enhanced I/O Type 8-Bit OTP MCU

HT48R016/HT48R017

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Table of Contents

Features	5
CPU Features	5
Peripheral Features.....	5
General Description	6
Selection Table	6
Block Diagram	6
Pin Assignment	6
Pin Description	7
Absolute Maximum Ratings	9
D.C.Characteristics	9
A.C. Characteristics	10
LVR Characteristics	11
Power-on Reset Characteristics	11
System Architecture	12
Clocking and Pipelining	12
Program Counter – PC.....	13
Stack	13
Arithmetic and Logic Unit – ALU	14
Program Memory	15
Structure.....	15
Special Vectors	16
Look-up Table.....	16
Data Memory	18
Structure.....	18
Special Purpose Data Memory	19
Special Function Register	20
Indirect Addressing Registers – IAR0, IAR1	20
Memory Pointers – MP0, MP1	20
Accumulator – ACC.....	21
Program Counter Low Register – PCL.....	21
Status Register – STATUS	21
Input/Output Ports and Control Registers	23
System Control Registers – CTRL0, CTRL1.....	23
Wake-up Function Register – PAWK.....	24
Pull-high Registers – PAPU, PBPU.....	24

Oscillator	25
System Oscillator Overview	25
System Clock Configurations	25
External Crystal/Resonator Oscillator – HXT	25
Internal RC Oscillator – HIRC	26
Internal Low Speed Oscillator – LIRC	26
Power Down Mode and Wake-up.....	27
Entering the Power Down Mode	27
Standby Current Considerations	27
Wake-up	28
Watchdog Timer.....	29
Watchdog Timer Clock Source.....	29
Watchdog Timer Control Register	29
Watchdog Timer Operation	30
Reset and Initialisation.....	32
Reset Functions	32
Reset Initial Conditions	35
Input/Output Ports	37
Pull-high Resistors	37
Port A Wake-up	37
I/O Port Control Registers	38
Pin-shared Functions	38
I/O Pin Structures.....	39
Programming Considerations.....	40
Timer/Event Counter	41
Configuring the Timer/Event Counter Input Clock Source	41
Timer Register – TMR.....	42
Timer Control Register – TMRC.....	42
Timer Mode	44
Event Counter Mode	44
Pulse Width Capture Mode	45
Prescaler.....	46
PFD Function	46
I/O Interfacing.....	48
Programming Considerations.....	48
Timer Program Example	49
Time Base.....	49

Interrupts	50
Interrupt Register	50
Interrupt Operation	51
Interrupt Priority	53
External Interrupt	53
Timer/Event Counter Interrupt	54
Time Base Interrupts	54
Interrupt Wake-up Function	54
Programming Considerations	55
Configuration Options	55
Application Circuits	55
Instruction Set	56
Introduction	56
Instruction Timing	56
Moving and Transferring Data	56
Arithmetic Operations	56
Logical and Rotate Operations	57
Branches and Control Transfer	57
Bit Operations	57
Table Read Operations	57
Other Operations	57
Instruction Set Summary	58
Table Conventions	58
Instruction Definition	60
Package Information	70
16-pin DIP (300mil) Outline Dimensions	70
16-pin NSOP (150mil) Outline Dimensions	73
Reel Dimensions	74
Carrier Tape Dimensions	75

Features

CPU Features

- Operating voltage:
 - ◆ $f_{SYS}=4\text{MHz}$: 2.2V~5.5V
 - ◆ $f_{SYS}=8\text{MHz}$: 3.3V~5.5V
 - ◆ $f_{SYS}=12\text{MHz}$: 4.5V~5.5V
- Up to 0.33 μs instruction cycle with 12MHz system clock at $V_{DD}=5\text{V}$
- Oscillator types:
 - ◆ External high frequency crystal -- HXT
 - ◆ Internal RC -- HIRC
- Power down modes and wake-up function
- Fully integrated internal 4MHz, 8MHz and 12MHz oscillator requires no external components
- LIRC oscillator function for watchdog timer
- All instructions executed in one or two instruction cycles
- Table read instructions
- 63 powerful instructions
- Up to 4-level stack
- Bit manipulation instruction

Peripheral Features

- Up to 2048 \times 16 program memory
- Up to 96 \times 8 data memory
- Up to 14 bidirectional I/O lines
- Low voltage reset function
- Watchdog timer function
- External interrupt input shared with an I/O line
- One 8-bit programmable Timer/Event Counter with overflow interrupt and prescaler
- Time-Base functions
- Programmable Frequency Divider – PFD
- Package Types: 16 DIP/NSOP

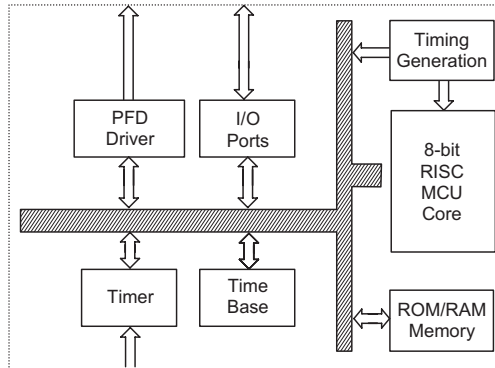
General Description

The Enhanced I/O MCUs are a series of 8-bit high performance, RISC architecture microcontrollers specifically designed for a wide range of applications. The usual Holtek microcontroller features of low power consumption, I/O flexibility, timer functions, oscillator options, power down and wake-up functions, watchdog timer and low voltage reset, combine to provide devices with a huge range of functional options while still maintaining a high level of cost effectiveness. The fully integrated system oscillator HIRC, which requires no external components and which has three frequency selections, opens up a huge range of new application possibilities for these devices, some of which may include industrial control, consumer products, household appliances subsystem controllers, etc.

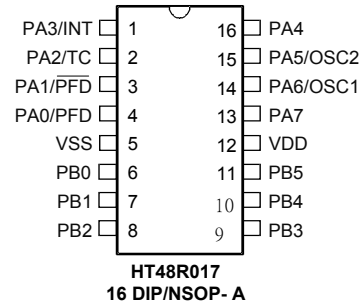
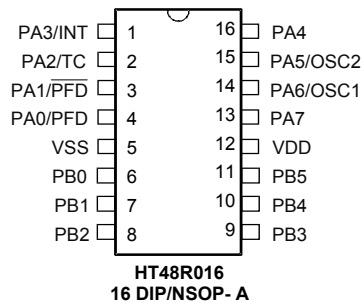
Selection Table

Part No.	ROM	RAM	I/O	8-bit Timer	Time Base	HIRC (MHz)	PFD	Stack	Package
HT48R016	1K×16	64×8	14	1	1	4/8/12	√	2	16DIP/NSOP
HT48R017	2K×16	96×8	14	1	1	4/8/12	√	4	16DIP/NSOP

Block Diagram



Pin Assignment



Pin Description

HT48R016

Pin Name	Function	OPT	I/T	O/T	Description
PA0/PFD	PA0	PAPU PAWK	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	PFD	CTRL0	—	CMOS	PFD output
PA1/ $\overline{\text{PFD}}$	PA1	PAPU PAWK	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	$\overline{\text{PFD}}$	CTRL0	—	CMOS	PFD Complementary Output
PA2/TC	PA2	PAPU PAWK	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	TC	TMRC	ST	—	External Timer 0 clock input
PA3/INT	PA3	PAPU PAWK	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	INT	CTRL0	ST	—	External interrupt input
PA4	PA4	PAPU PAWK	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
PA5/OSC2	PA5	PAPU PAWK	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	OSC2	CO	—	HXT	Oscillator pin
PA6/OSC1	PA6	PAPU PAWK	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	OSC1	CO	HXT	—	Oscillator pin
PA7	PA7	PAWK	ST	NMOS	General purpose I/O. Register enabled wake-up
PB0~PB5	PBn	PBPU	ST	CMOS	General purpose I/O. Register enabled pull-up
VDD	VDD	—	PWR	—	Power supply
VSS	VSS	—	PWR	—	Ground

HT48R017

Pin Name	Function	OPT	I/T	O/T	Description
PA0/PFD	PA0	PAPU PAWK	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	PFD	CTRL0	—	CMOS	PFD output
PA1/PFD	PA1	PAPU PAWK	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	PFD	CTRL0	—	CMOS	PFD Complementary Output
PA2/TC	PA2	PAPU PAWK	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	TC	TMRC	ST	—	External Timer 0 clock input
PA3/INT	PA3	PAPU PAWK	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	INT	CTRL0	ST	—	External interrupt input
PA4	PA4	PAPU PAWK	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
PA5/OSC2	PA5	PAPU PAWK	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	OSC2	CO	—	HXT	Oscillator pin
PA6/OSC1	PA6	PAPU PAWK	ST	CMOS	General purpose I/O. Register enabled pull-up and wake-up
	OSC1	CO	HXT	—	Oscillator pin
PA7	PA7	PAWK	ST	NMOS	General purpose I/O. Register enabled wake-up
PB0~PB5	PBn	PBPU	ST	CMOS	General purpose I/O. Register enabled pull-up
VDD	VDD	—	PWR	—	Power supply
VSS	VSS	—	PWR	—	Ground

Note: I/T: Input type; O/T: Output type
 OPT: Optional by configuration option (CO) or register option
 PWR: Power; CO: Configuration option
 ST: Schmitt Trigger input; CMOS: CMOS output; AN: Analog input
 HXT: High frequency crystal oscillator

Absolute Maximum Ratings

Supply Voltage	$V_{SS}-0.3V$ to $V_{SS}+6.0V$
Input Voltage	$V_{SS}-0.3V$ to $V_{DD}+0.3V$
Storage Temperature.....	$-50^{\circ}C$ to $125^{\circ}C$
Operating Temperature.....	$-40^{\circ}C$ to $85^{\circ}C$
I_{OH} Total	-100mA
I_{OL} Total	100mA
Total Power Dissipation	500mW

Note: These are stress ratings only. Stresses exceeding the range specified under "Absolute Maximum Ratings" may cause substantial damage to the device. Functional operation of this device at other conditions beyond those listed in the specification is not implied and prolonged exposure to extreme conditions may affect device reliability.

D.C.Characteristics

$T_a=25^{\circ}C$

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V_{DD}	Conditions				
V_{DD}	Operating Voltage	—	$f_{SYS}=4MHz$	2.2	—	5.5	V
			$f_{SYS}=8MHz$	3.3	—	5.5	V
			$f_{SYS}=12MHz$	4.5	—	5.5	V
I_{DD1}	Operating Current (HXT, HIRC)	3V	No load, $f_{SYS}=4MHz$	—	1	2	mA
		5V		—	2.5	5	mA
I_{DD2}	Operating Current (HXT, HIRC)	5V	No load, $f_{SYS}=8MHz$	—	4	8	mA
I_{DD3}	Operating Current (HXT, HIRC)	5V	No load, $f_{SYS}=12MHz$	—	6	12	mA
I_{STB1}	Standby Current (LIRC on)	3V	No load, system HALT	—	—	5	μA
		5V		—	—	10	μA
I_{STB2}	Standby Current (LIRC off)	3V	No load, system HALT	—	—	1	μA
		5V		—	—	2	μA
V_{IL1}	Input Low Voltage for PA, PB, TC, INT (except PA7)	5V	—	0	—	1.5	V
		—	—	0	—	$0.2V_{DD}$	V
V_{IH1}	Input High Voltage for PA, PB, TC, INT (except PA7)	5V	—	3.5	—	5	V
		—	—	$0.8V_{DD}$	—	V_{DD}	V
V_{IL2}	Input Low Voltage (PA7)	—	—	0	—	$0.4V_{DD}$	V
V_{IH2}	Input High Voltage (PA7)	—	—	$0.9V_{DD}$	—	V_{DD}	V
I_{OL1}	I/O Sink Current (PA,PB)	3V	$V_{OL}=0.1V_{DD}$	4	8	—	mA
		5V		10	20	—	mA
I_{OH1}	I/O Source Current (PA,PB)	3V	$V_{OH}=0.9V_{DD}$	-2	-4	—	mA
		5V		-5	-10	—	mA
I_{OL2}	PA7 Sink Current	5V	$V_{OL}=0.1V_{DD}$	2	3	—	mA
R_{PH}	Pull-high Resistance (I/O)	3V	—	20	60	100	k Ω
		5V	—	10	30	50	k Ω

Note: The standby current (I_{STB1} ~ I_{STB2}) are measured with all I/O pins in input mode and tied to VDD.

A.C. Characteristics

Ta=25°C

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V _{DD}	Condition				
f _{SYS}	System clock	—	2.2~5.5V	32	—	4000	kHz
			3.3~5.5V	32	—	8000	
			4.5~5.5V	32	—	12000	
f _{HIRC}	System clock (HIRC)	3.3V/5V	Ta=25°C	-2%	4/8	+2%	MHz
		5V	Ta=25°C	-2%	12	+2%	
		3.3V/5V	Ta=0~70°C	-5%	4/8	+5%	
		5V	Ta=0~70°C	-5%	12	+5%	
		2.2V~3.6V	Ta=0~70°C	-8%	4	+8%	
		3.3V~5.5V	Ta=0~70°C	-8%	4/8	+8%	
		4.5V~5.5V	Ta=0~70°C	-8%	12	+8%	
		2.2V~3.6V	Ta=-40°C~85°C	-12%	4	+12%	
		3.3V~5.5V	Ta=-40°C~85°C	-12%	4/8	+12%	
4.5V~5.5V	Ta=-40°C~85°C	-12%	12	+12%			
f _{LIRC}	System Clock (LIRC)	5V	—	28.8	32	35.2	kHz
		2.2V~5.5V	Ta=-40°C~85°C	16	32	51.2	
t _{SST}	System start-up timer period (wake-up from power down mode)	—	f _{sys} =HXT	1024	—	—	t _{sys}
			f _{sys} =HIRC	2	—	—	
t _{TC}	TC input pin pulse width	—	—	0.3	—	—	μs
t _{INT}	Interrupt input pin pulse width	—	—	10	—	—	μs
t _{RSTD}	System Reset Delay Time (Power on reset)	—	—	25	50	100	ms
	System Reset Delay Time (Any reset except Power on reset)	—	—	8.3	16.7	33.3	ms

 Note: 1. t_{sys}=1/f_{sys}

- To maintain the accuracy of the internal HIRC oscillator frequency, a 0.1F decoupling capacitor should be connected between VDD and VSS and located as close to the device as possible.

LVR Characteristics

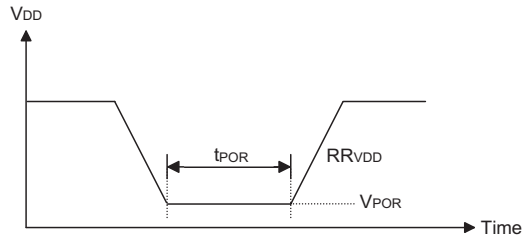
Ta=25°C

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V _{DD}	Conditions				
V _{LVR1}	Low Voltage Reset Voltage	—	LVR Enable, 2.1V selected	-5% Typ.	2.1	+5% Typ.	V
V _{LVR2}			LVR Enable, 2.55V selected		2.55		V
V _{LVR3}			LVR Enable, 3.15V selected		3.15		V
V _{LVR4}			LVR Enable, 3.8V selected		3.8		V
I _{LVR}	Additional Power Consumption	3V	LVR enabled	—	30	45	μA
		5V		—	60	90	μA
t _{LVR}	Low Voltage Width to Reset	—	—	120	240	480	μs
t _{SRESET}	Software Reset Width to Reset	—	—	45	90	120	μs

Power-on Reset Characteristics

Ta=25°C

Symbol	Parameter	Test Conditions		Min.	Typ.	Max.	Unit
		V _{DD}	Conditions				
V _{POR}	V _{DD} Start Voltage to Ensure Power-on Reset	—	—	—	—	100	mV
RR _{VDD}	V _{DD} Raising Rate to Ensure Power-on Reset	—	—	0.035	—	—	V/ms
t _{POR}	Minimum Time for V _{DD} Stays at V _{POR} to Ensure Power-on Reset	—	—	1	—	—	ms



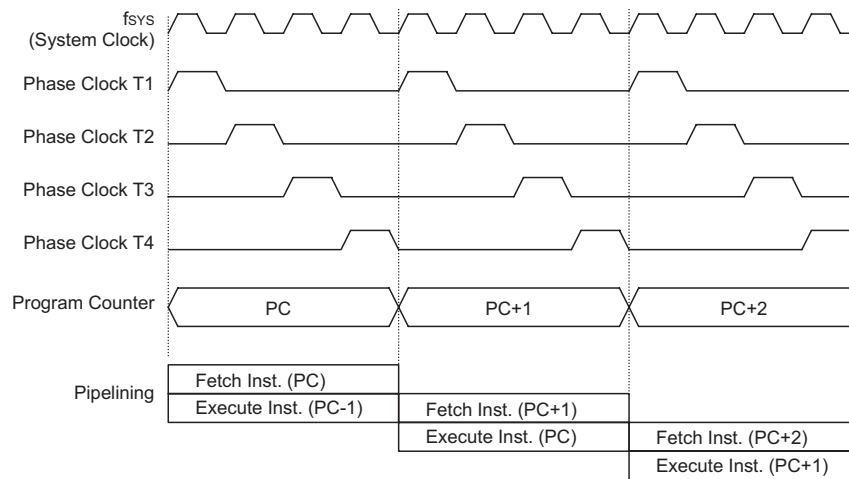
System Architecture

A key factor in the high-performance features of the Holtek range of microcontrollers is attributed to their internal system architecture. The range of device takes advantage of the usual features found within RISC microcontrollers providing increased speed of operation and enhanced performance. The pipelining scheme is implemented in such a way that instruction fetching and instruction execution are overlapped, hence instructions are effectively executed in one cycle, with the exception of branch or call instructions. An 8-bit wide ALU is used in practically all instruction set operations, which carries out arithmetic operations, logic operations, rotation, increment, decrement, branch decisions, etc. The internal data path is simplified by moving data through the Accumulator and the ALU. Certain internal registers are implemented in the Data Memory and can be directly or indirectly addressed. The simple addressing methods of these registers along with additional architectural features ensure that a minimum of external components is required to provide a functional I/O control system with maximum reliability and flexibility.

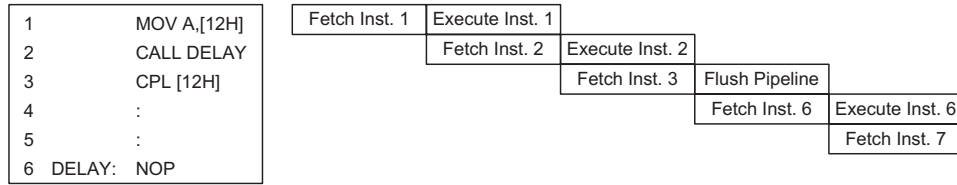
Clocking and Pipelining

The main system clock, derived from either a Crystal/Resonator or RC oscillator is subdivided into four internally generated non-overlapping clocks, T1~T4. The Program Counter is incremented at the beginning of the T1 clock during which time a new instruction is fetched. The remaining T2~T4 clocks carry out the decoding and execution functions. In this way, one T1~T4 clock cycle forms one instruction cycle. Although the fetching and execution of instructions takes place in consecutive instruction cycles, the pipelining structure of the microcontroller ensures that instructions are effectively executed in one instruction cycle. The exception to this are instructions where the contents of the Program Counter are changed, such as subroutine calls or jumps, in which case the instruction will take one more instruction cycle to execute.

For instructions involving branches, such as jump or call instructions, two instruction cycles are required to complete instruction execution. An extra cycle is required as the program takes one cycle to firstly obtain the actual jump or call address and then another cycle to actually execute the branch. The requirement for this extra cycle should be taken into account by programmers in timing sensitive applications.



System Clocking and Pipelining



Instruction Fetching

Program Counter – PC

During program execution, the Program Counter is used to keep track of the address of the next instruction to be executed. It is automatically incremented by one each time an instruction is executed except for instructions, such as “JMP” or “CALL” that demand a jump to a non-consecutive Program Memory address. It must be noted that only the lower 8 bits, known as the Program Counter Low Register, are directly addressable by user.

When executing instructions requiring jumping to non-consecutive addresses such as a jump instruction, a subroutine call, interrupt or reset, etc, the microcontroller manages program control by loading the required address into the Program Counter. For conditional skip instructions, once the condition has been met, the next instruction, which has already been fetched during the present instruction execution, is discarded and a dummy cycle takes its place while the correct instruction is obtained.

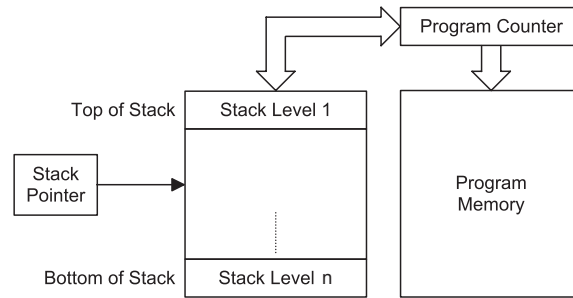
Device	Program Counter	
	Program Counter High Byte	PCL Register
HT48R016	PC9, PC8	PCL7~PCL0
HT48R017	PC10~PC8	

The lower byte of the Program Counter, known as the Program Counter Low register or PCL, is available for program control and is a readable and writeable register. By transferring data directly into this register, a short program jump can be executed directly. However, as only this low byte is available for manipulation, the jumps are limited in the present page of memory, which have 256 locations. When such program jumps are executed it should also be noted that a dummy cycle will be inserted.

The lower byte of the Program Counter is fully accessible under program control. Manipulating the PCL might cause program branching, so an extra cycle is needed to pre-fetch. Further information on the PCL register can be found in the Special Function Register section.

Stack

This is a special part of the memory which is used to save the contents of the Program Counter only. The stack is neither part of the data nor part of the program space, and is neither readable nor writeable. The activated level is indexed by the Stack Pointer, and is neither readable nor writeable. At a subroutine call or interrupt acknowledge signal, the contents of the Program Counter are pushed onto the stack. At the end of a subroutine or an interrupt routine, signaled by a return instruction, RET or RETI, the Program Counter is restored to its previous value from the stack. After a device reset, the Stack Pointer will point to the top of the stack.



Device	Stack Levels
HT48R016	2
HT48R017	4

If the stack is full and an enabled interrupt takes place, the interrupt request flag will be recorded but the acknowledge signal will be inhibited. When the Stack Pointer is decremented, by RET or RETI, the interrupt will be serviced. This feature prevents stack overflow allowing the programmer to use the structure more easily. However, when the stack is full, a CALL subroutine instruction can still be executed which will result in a stack overflow. Precautions should be taken to avoid such cases which might cause unpredictable program branching.

Arithmetic and Logic Unit – ALU

The arithmetic-logic unit or ALU is a critical area of the microcontroller that carries out arithmetic and logic operations of the instruction set. Connected to the main microcontroller data bus, the ALU receives related instruction codes and performs the required arithmetic or logical operations after which the result will be placed in the specified register. As these ALU calculation or operations may result in carry, borrow or other status changes, the status register will be correspondingly updated to reflect these changes. The ALU supports the following functions:

- Arithmetic operations: ADD, ADDM, ADC, ADCM, SUB, SUBM, SBC, SBCM, DAA
- Logic operations: AND, OR, XOR, ANDM, ORM, XORM, CPL, CPLA
- Rotation RRA, RR, RRCA, RRC, RLA, RL, RLCA, RLC
- Increment and Decrement INCA, INC, DECA, DEC
- Branch decision, JMP, SZ, SZA, SNZ, SIZ, SDZ, SIZA, SDZA, CALL, RET, RETI.

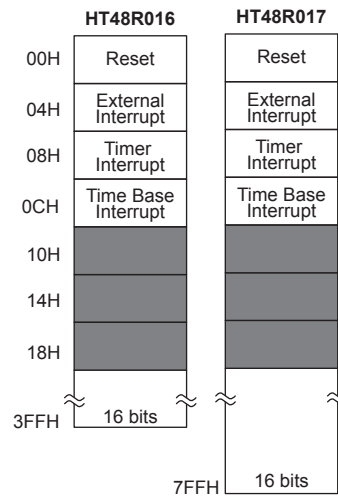
Program Memory

The Program Memory is the location where the user code or program is stored. The device is supplied with One-Time Programmable, OTP, memory where users can program their application code into the device. By using the appropriate programming tools, OTP device offers users the flexibility to freely develop their applications which may be useful during debug or for products requiring frequent upgrades or program changes.

Structure

The Program Memory has a capacity of 1k×16 to 2k×16. The Program Memory is addressed by the Program Counter and also contains data, table information and interrupt entries information. Table data, which can be setup in any location within the Program Memory, is addressed by separate table pointer registers.

Device	Capacity
HT48R016	1k×16
HT48R017	2k×16



Program Memory Structure

Special Vectors

Within the Program Memory, certain locations are reserved for special usage such as reset and interrupts.

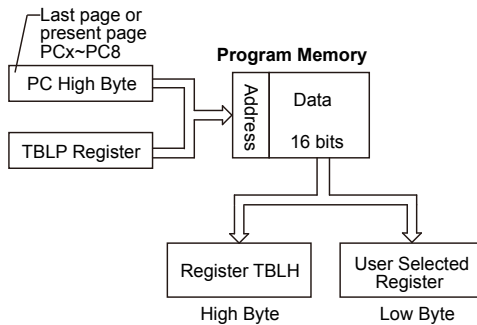
- **Reset Vector**
 This vector is reserved for use by the device reset for program initialisation. After a device reset is initiated, the program will jump to this location and begin execution.
- **External interrupt vector**
 This vector is used by the external interrupt. If the external interrupt pin on the device receives an edge transition, the program will jump to this location and begin execution if the external interrupt is enabled and the stack is not full. The external interrupt active edge transition type, whether high to low, low to high or both is specified in the CTRL0 register.
- **Timer/Event counter interrupt vector**
 This internal vector is used by the Timer/Event Counters. If a Timer/Event Counter overflow occurs, the program will jump to its respective location and begin execution if the associated Timer/Event Counter interrupt is enabled and the stack is not full.
- **Time base interrupt vector**
 This internal vector is used by the internal Time Base. If a Time Base overflow occurs, the program will jump to this location and begin execution if the Time Base counter interrupt is enabled and the stack is not full.

Look-up Table

Any location within the Program Memory can be defined as a look-up table where programmers can store fixed data. To use the look-up table, the table pointer must first be setup by placing the address of the look up data to be retrieved in the table pointer register, TBLP. These registers define the total address of the look-up table.

After setting up the table pointer, the table data can be retrieved from the current Program Memory page or last Program Memory page using the “TABRDC [m]” or “TABRDL [m]” instructions, respectively. When these instructions are executed, the lower order table byte from the Program Memory will be transferred to the user defined Data Memory register [m] as specified in the instruction. The higher order table data byte from the Program Memory will be transferred to the TBLH special register. Any unused bits in this transferred higher order byte will be read as “0”.

The following diagram illustrates the addressing data flow of the look-up table:



Instruction	Table Location Bits										
	b10	b9	b8	b7	b6	b5	b4	b3	b2	b1	b0
TABRDC [m]	PC10	PC9	PC8	@7	@6	@5	@4	@3	@2	@1	@0
TABRDL [m]	1	1	1	@7	@6	@5	@4	@3	@2	@1	@0

Table Location

Note: PC10~PC8: Current Program Counter bits

@7~@0: Table Pointer TBLP bits

For the HT48R016, the Table address location is 10 bits, i.e. from b9~b0.

For the HT48R017, the Table address location is 11 bits, i.e. from b10~b0.

Table Read Program Example:

```

tempreg1 db ?      ; temporary register #1
tempreg2 db ?      ; temporary register #2
:
:
mov a,06h          ; initialise table pointer - note that this address
                  ; is referenced to the last page or present page

mov tblp, a
:
:
tabrdl tempreg1   ; transfers value in table referenced by table pointer
                  ; to tempreg1 data at prog.memory address "306H"
                  ; transferred to to tempreg1 and TBLH

dec tblp          ; reduce value of table pointer by one
tabrdl tempreg2   ; transfers value in table referenced by table pointer
                  ; to tempreg2 data at prog.memory address "305H"
                  ; transferred to tempreg2 and TBLH in this example the
                  ; data "1AH" is transferred to tempreg1 and data "0FH"
                  ; to register tempreg2 the value "00H" will be
                  ; transferred to the high byte register TBLH

:
:
org 300h          ; sets initial address of last page
dc 00Ah, 00Bh, 00Ch, 00Dh, 00Eh, 00Fh, 01Ah, 01Bh
:
:

```

Data Memory

The Data Memory is a volatile area 8-bit wide RAM internal memory and is the location where temporary information is stored.

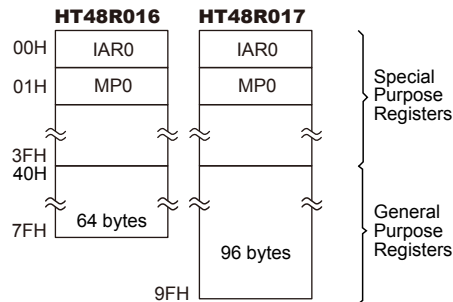
Structure

Divided into two sections, the first of these is an area of RAM where special function registers are located. These registers have fixed locations and are necessary for correct operation of the device. Many of these registers can be read from and written to directly under program control, however, some remain protected from user manipulation. The second area of Data Memory is reserved for general purpose use. All locations within this area are read and write accessible under program control.

Device	Capacity
HT48R016	64×8
HT48R017	96×8

The two sections of Data Memory, the Special Purpose and General Purpose Data Memory are located at consecutive locations. All are implemented in RAM and are 8 bits wide but the length of each memory section is dictated by the type of microcontroller chosen. The start address of the Data Memory for all devices is the address "00H".

All microcontroller programs require an area of read/write memory where temporary data can be stored and retrieved for use later. It is this area of RAM memory that is known as General Purpose Data Memory. This area of Data Memory is fully accessible by the user program for both read and write operations. By using the "SET [m].i" and "CLR [m].i" instructions individual bits can be set or reset under program control giving the user a large range of flexibility for bit manipulation in the Data Memory.



Data Memory Structure

Note: Most of the Data Memory bits can be directly manipulated using the "SET [m].i" and "CLR [m].i" with the exception of a few dedicated bits. The Data Memory can also be accessed through the memory pointer registers.

Special Function Register

To ensure successful operation of the microcontroller, certain internal registers are implemented in the Data Memory area. These registers ensure correct operation of internal functions such as timers, interrupts, etc., as well as external functions such as I/O data control. The location of these registers within the Data Memory begins at the address "00H". Any unused Data Memory locations between these special function registers and the point where the General Purpose Memory begins is reserved and attempting to read data from these locations will return a value of "00H".

Indirect Addressing Registers – IAR0, IAR1

The Indirect Addressing Registers, IAR0 and IAR1, although having their locations in normal RAM register space, do not actually physically exist as normal registers. The method of indirect addressing for RAM data manipulation uses these Indirect Addressing Registers and Memory Pointers, in contrast to direct memory addressing, where the actual memory address is specified. Actions on the IAR0 and IAR1 registers will result in no actual read or write operation to these registers but rather to the memory location specified by their corresponding Memory Pointer, MP0 or MP1. Acting as a pair, IAR0 with MP0 and IAR1 with MP1 can together access data from the Data Memory. As the Indirect Addressing Registers are not physically implemented, reading the Indirect Addressing Registers indirectly will return a result of "00H" and writing to the registers indirectly will result in no operation.

Memory Pointers – MP0, MP1

Two Memory Pointers, known as MP0 and MP1 are provided. These Memory Pointers are physically implemented in the Data Memory and can be manipulated in the same way as normal registers providing a convenient way with which to indirectly address and track data. When any operation to the relevant Indirect Addressing Registers is carried out, the actual address that the microcontroller is directed to, is the address specified by the related Memory Pointer. The following example shows how to clear a section of four Data Memory locations already defined as locations adres1 to adres4.

Indirect Addressing Program Example

```
data . section 'data'
adres1    db    ?
adres2    db    ?
adres3    db    ?
adres4    db    ?
block     db    ?
code. section at 0 code
org 00h
start:
mov a,04h                ; setup size of block
mov block,a
mov a,offset adres1      ; Accumulator loaded with first RAM address
mov mp0,a                ; setup memory pointer with first RAM address
loop:
clr IAR0                 ; clear the data at address defined by MP0
inc mp0                   ; increment memory pointer
sdz block                 ; check if last memory location has been cleared
jmp loop
continue:
```

The important point to note here is that in the example shown above, no reference is made to specific Data Memory addresses.

Accumulator – ACC

The Accumulator is central to the operation of any microcontroller and is closely related with operations carried out by the ALU. The Accumulator is the place where all intermediate results from the ALU are stored. Without the Accumulator it would be necessary to write the result of each calculation or logical operation such as addition, subtraction, shift, etc., to the Data Memory resulting in higher programming and timing overheads. Data transfer operations usually involve the temporary storage function of the Accumulator; for example, when transferring data between one user-defined register and another, it is necessary to do this by passing the data through the Accumulator as no direct transfer between two registers is permitted.

Program Counter Low Register – PCL

To provide additional program control functions, the low byte of the Program Counter is made accessible to programmers by locating it within the Special Purpose area of the Data Memory. By manipulating this register, direct jumps to other program locations are easily implemented. Loading a value directly into this PCL register will cause a jump to the specified Program Memory location, however as the register is only 8-bit wide only jumps within the current Program Memory page are permitted. When such operations are used, note that a dummy cycle will be inserted.

Status Register – STATUS

This 8-bit register contains the zero flag (Z), carry flag (C), auxiliary carry flag (AC), overflow flag (OV), power down flag (PDF), and watchdog time-out flag (TO). These arithmetic/logical operation and system management flags are used to record the status and operation of the microcontroller.

With the exception of the TO and PDF flags, bits in the status register can be altered by instructions like most other registers. Any data written into the status register will not change the TO or PDF flag. In addition, operations related to the status register may give different results due to the different instruction operations. The TO flag can be affected only by a system power-up, a WDT time-out or by executing the “CLR WDT” or “HALT” instruction. The PDF flag is affected only by executing the “HALT” or “CLR WDT” instruction or during a system power-up.

The Z, OV, AC and C flags generally reflect the status of the latest operations.

In addition, on entering an interrupt sequence or executing a subroutine call, the status register will not be pushed onto the stack automatically. If the contents of the status registers are important and if the subroutine can corrupt the status register, precautions must be taken to correctly save it. Note that bits 0~3 of the STATUS register are both readable and writeable bits.

STATUS Register

Bit	7	6	5	4	3	2	1	0
Name	—	—	TO	PDF	OV	Z	AC	C
R/W	—	—	R/W	R/W	R/W	R/W	R/W	R/W
POR	—	—	0	0	x	x	x	x

“x” unknown

- Bit 7~6 Unimplemented, read as “0”
- Bit 5 **TO:** Watchdog Time-Out flag
 0: after power up or executing the “CLR WDT” or “HALT” instruction
 1: a watchdog time-out occurred
- Bit 4 **PDF:** Power down flag
 0: after power up or executing the “CLR WDT” instruction
 1: by executing the “HALT” instruction
- Bit 3 **OV:** Overflow flag
 0: no overflow
 1: an operation results in a carry into the highest-order bit but not a carry out of the
 highest-order bit or vice versa
- Bit 2 **Z:** Zero flag
 0: the result of an arithmetic or logical operation is not zero
 1: the result of an arithmetic or logical operation is zero
- Bit 1 **AC:** Auxiliary flag
 0: no auxiliary carry
 1: an operation results in a carry out of the low nibbles in addition, or no borrow
 from the high nibble into the low nibble in subtraction
- Bit 0 **C:** Carry flag
 0: no carry-out
 1: an operation results in a carry during an addition operation or if a borrowing does
 not take place during a subtraction operation
 C is also affected by a rotate through carry instruction

Input/Output Ports and Control Registers

Within the area of Special Function Registers, the port PA, PB, etc data I/O registers and their associated control register PAC, PBC, etc play a prominent role. These registers are mapped to specific addresses within the Data Memory as shown in the Data Memory table. The data I/O registers, are used to transfer the appropriate output or input data on the port. The control registers specifies which pins of the port are set as inputs and which are set as outputs. To setup a pin as an input, the corresponding bit of the control register must be set high, for an output it must be set low. During program initialisation, it is important to first setup the control registers to specify which pins are outputs and which are inputs before reading data from or writing data to the I/O ports. One flexible feature of these registers is the ability to directly program single bits using the "SET [m].i" and "CLR [m].i" instructions. The ability to change I/O pins from output to input and vice versa by manipulating specific bits of the I/O control registers during normal program operation is a useful feature of these devices.

System Control Registers – CTRL0, CTRL1

These registers are used to provide control over various internal functions. Some of these include the PFD control, certain system clock options, external Interrupt edge trigger type, Watchdog Timer enable function, Time Base function division ratio.

CTRL0 Register

Bit	7	6	5	4	3	2	1	0
Name	INTEG1	INTEG0	TBSEL1	TBSEL0	—	—	—	—
R/W	R/W	R/W	R/W	R/W	—	—	—	—
POR	1	0	0	0	—	—	—	—

- Bit 7~6 **INTEG1, INTEG0:** External interrupt edge type
 00: disable
 01: rising edge trigger
 10: falling edge trigger
 11: dual edge trigger
- Bit 5~4 **TBSEL1, TBSEL0:** Time base period selection
 00: $2^{10}/f_{TP}$
 01: $2^{11}/f_{TP}$
 10: $2^{12}/f_{TP}$
 11: $2^{13}/f_{TP}$
- Bit 3~0 Unimplemented, read as "0"

CTRL1 Register

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	LVRF	LRF	WRF
R/W	—	—	—	—	—	R/W	R/W	R/W
POR	—	—	—	—	—	x	0	0

“x” unknown

- Bit 7~3 Unimplemented, read as "0"
- Bit 2 **LVRF:** LVR function reset flag
 0: Not occur
 1: Occurred
 This bit is set to 1 when a specific Low Voltage Reset situation condition occurs. This bit can only be cleared to 0 by the application program.
- Bit 1 **LRF:** LVR Control register software reset flag
 0: Not occur
 1: Occurred
 This bit is set to 1 if the LVRC register contains any non defined LVR voltage register values. This in effect acts like a software reset function. This bit can only be cleared to 0 by the application program.
- Bit 0 **WRF:** WDT Control register software reset flag
 0: Not occur
 1: Occurred
 This bit is set to 1 by the WDT Control register software reset and cleared by the application program. Note that this bit can only be cleared to 0 by the application program.

Wake-up Function Register – PAWK

When the microcontroller enters the Sleep Mode, various methods exist to wake the device up and continue with normal operation. One method is to allow a falling edge on the I/O pins to have a wake-up function. This register is used to select which Port A I/O pins are used to have this wake-up function.

Pull-high Registers – PAPU, PBPU

The I/O pins, if configured as inputs, can have internal pull-high resistors connected, which eliminates the need for external pull-high resistors. This register selects which I/O pins are connected to internal pull-high resistors.

Oscillator

Various oscillator options offer the user a wide range of functions according to their various application requirements. The flexible features of the oscillator functions ensure that the best optimisation can be achieved in terms of speed and power saving. Oscillator selections and operation are selected through a combination of configuration options and registers.

System Oscillator Overview

In addition to being the source of the main system clock the oscillators also provide clock sources for the Watchdog Timer and Time Base functions. External oscillators requiring some external components as well as a two fully integrated internal oscillators, requiring no external components, are provided to form a wide range of both fast and slow system oscillators.

Type	Name	Freq.	Pins
External Crystal	HXT	400kHz~12MHz	OSC1/OSC2
Internal High Speed RC	HIRC	4, 8 or 12MHz	—
Internal Low Speed RC	LIRC	32kHz	—

Oscillator Types

System Clock Configurations

There are two system oscillators. These two system oscillators are the external crystal/ceramic oscillator – HXT, and the internal RC oscillator – HIRC. The low speed oscillator is the internal 32kHz ($V_{DD}=5V$) oscillator – LIRC.

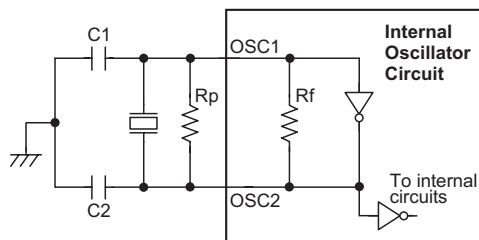
External Crystal/Resonator Oscillator – HXT

The simple connection of a crystal across OSC1 and OSC2 will create the necessary phase shift and feedback for oscillation. However, for some crystals and most resonator types, to ensure oscillation and accurate frequency generation, it is necessary to add two small value external capacitors, C1 and C2. The exact values of C1 and C2 should be selected in consultation with the crystal or resonator manufacturer's specification.

Crystal Oscillator C1 and C2 Values		
Crystal Frequency	C1	C2
12MHz	8pF	10pF
8MHz	8pF	10pF
4MHz	8pF	10pF
1MHz	100pF	100pF

Note: C1 and C2 values are for guidance only.

Crystal Recommended Capacitor Values

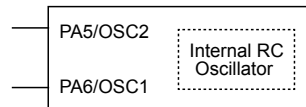


- Note: 1. R_p is normally not required. C1 and C2 are required.
 2. Although not shown OSC1/OSC2 pins have a parasitic capacitance of around 7pF.

Crystal/Resonator Oscillator – HXT

Internal RC Oscillator – HIRC

The internal RC oscillator is a fully integrated system oscillator requiring no external components. The internal RC oscillator has three fixed frequencies of either 4MHz, 8MHz or 12MHz. Device trimming during the manufacturing process and the inclusion of internal frequency compensation circuits are used to ensure that the influence of the power supply voltage, temperature and process variations on the oscillation frequency are minimised. As a result, at a power supply of either 3V or 5V and at a temperature of 25 degrees, the fixed oscillation frequency of 4MHz, 8MHz or 12MHz will have a tolerance within 2%. Note that if this internal system clock option is selected, as it requires no external pins for its operation, I/O pins PA5 and PA6 are free for use as normal I/O pins.



Note: PA5/PA6 used as normal I/Os

Internal RC Oscillator – HIRC

Internal Low Speed Oscillator – LIRC

The LIRC is a fully self-contained free running on-chip RC oscillator with a typical frequency of 32kHz at 5V requiring no external components. When the device enters the power down Mode, the system clock will stop running but the WDT oscillator continues to free-run and to keep the watchdog active.

Power Down Mode and Wake-up

Entering the Power Down Mode

There is only one way for the device to enter the Power Down Mode and that is to execute the “HALT” instruction in the application program. When this instruction is executed, the following will occur:

- The system oscillator will stop running and the application program will stop at the “HALT” instruction.
- The Data Memory contents and registers will maintain their present condition.
- The WDT will be cleared and resume counting if the WDT function is enabled.
- The I/O ports will maintain their present condition.
- In the status register, the Power Down flag, PDF, will be set and the Watchdog time-out flag, TO, will be cleared.

Standby Current Considerations

As the main reason for entering the Power Down Mode is to keep the current consumption of the MCU to as low a value as possible, perhaps only in the order of several micro-amps, there are other considerations which must also be taken into account by the circuit designer if the power consumption is to be minimised.

Special attention must be made to the I/O pins on the device. All high-impedance input pins must be connected to either a fixed high or low level as any floating input pins could create internal oscillations and result in increased current consumption. These must either be setup as outputs or if setup as inputs must have pull-high resistors connected. Care must also be taken with the loads, which are connected to I/O pins, which are setup as outputs. These should be placed in a condition in which minimum current is drawn or connected only to external circuits that do not draw current, such as other CMOS inputs.

The Watchdog Timer will continue to run when in the Power Down Mode and thus will consume some power since the Watchdog Timer is enabled and the clock source is derived from the LIRC oscillator.

Wake-up

After the system enters the Power Down Mode, it can be woken up from one of various sources listed as follows:

- An external falling edge on PA0~PA7
- A system interrupt
- A WDT overflow

If the system is woken up by Power-on reset, the device will experience a full system reset, however, if the device is woken up by a WDT overflow, a Watchdog Timer reset will be initiated. Although both of these wake-up methods will initiate a reset operation, the actual source of the wake-up can be determined by examining the TO and PDF flags. The PDF flag is cleared by a system power-up or executing the clear Watchdog Timer instructions and is set when executing the “HALT” instruction. The TO flag is set if a WDT time-out occurs, and causes a wake-up that only resets the Program Counter and Stack Pointer, the other flags remain in their original status.

Pins PA0 to PA7 can be setup via the PAWK register to permit a negative transition on the pin to wake-up the system. When a PA0 to PA7 pin wake-up occurs, the program will resume execution at the instruction following the “HALT” instruction.

If the system is woken up by an interrupt, then two possible situations may occur. The first is where the related interrupt is disabled or the interrupt is enabled but the stack is full, in which case the program will resume execution at the instruction following the “HALT” instruction. In this situation, the interrupt which wake-up the device will not be immediately serviced, but will rather be serviced later when the related interrupt is finally enabled or when a stack level becomes free. The other situation is where the related interrupt is enabled and the stack is not full, in which case the regular interrupt response takes place. If an interrupt request flag is set to “1” before entering the Power Down Mode, then any future interrupt requests will not generate a wake-up function of the related interrupt will be ignored.

No matter what the source of the wake-up event is, once a wake-up event occurs, there will be a time delay before normal program execution resumes. Consult the table for the related time.

Wake-up Source	Oscillator Type	
	HIRC	HXT
Power-on/LVR Reset	$t_{RSDT} + t_{SST2}$	$t_{RSDT} + t_{SST2}$
PA0~PA7 I/O port	t_{SST1}	t_{SST2}
Interrupt		
WDT Overflow		

Wake-up Delay Time

- Note: 1. t_{RSTD} (reset delay time), t_{SYS} (system clock)
 2. t_{RSTD} is power-on delay, typical time=50ms
 3. $t_{SST1} = 2t_{SYS}$
 4. $t_{SST2} = 1024t_{SYS}$

Watchdog Timer

The Watchdog Timer is provided to prevent program malfunctions or sequences from jumping to unknown locations, due to certain uncontrollable external events such as electrical noise.

Watchdog Timer Clock Source

The Watchdog Timer clock source is provided by the internal clock, f_s , which is in turn supplied by the LIRC oscillator. The LIRC internal oscillator has an approximate period of 32kHz at a supply voltage of 5V. However, it should be noted that this specified internal clock period can vary with V_{DD} , temperature and process variations. The Watchdog Timer source clock is then subdivided by a ratio of 2^8 to 2^{18} to give longer timeouts, the actual value being chosen using the WS2~WS0 bits in the WDTC register.

Watchdog Timer Control Register

A single register, WDTC, controls the required timeout period as well as the enable/disable operation. This register together with the corresponding configuration option control the overall operation of the Watchdog Timer.

WDTC Register

Bit	7	6	5	4	3	2	1	0
Name	WE4	WE3	WE2	WE1	WE0	WS2	WS1	WS0
R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	0	1	0	1	0	0	1	1

Bit 7~3 **WE4~WE0:** WDT function software control

10101: Disabled
 01010: Enabled
 Other: Reset MCU

When these bits are changed by the environmental noise to reset the microcontroller, the reset operation will be activated after 2~3 LIRC clock cycles and the WRF bit in the CTRL1 register will be set to 1.

Bit 2~0 **WS2~WS0:** WDT Time-out period selection

000: $2^8/f_s$
 001: $2^{10}/f_s$
 010: $2^{12}/f_s$
 011: $2^{14}/f_s$
 100: $2^{15}/f_s$
 101: $2^{16}/f_s$
 110: $2^{17}/f_s$
 111: $2^{18}/f_s$

CTRL1 Register

Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	LVRF	LRF	WRF
R/W	—	—	—	—	—	R/W	R/W	R/W
POR	—	—	—	—	—	x	0	0

“x” unknown

Bit 7~3 Unimplemented, read as "0"

Bit 2 **LVRF**: LVR function reset flag
Describe elsewhere.

Bit 1 **LRF**: LVR Control register software reset flag
Describe elsewhere.

Bit 0 **WRF**: WDT Control register software reset flag
0: Not occur
1: Occurred

This bit is set to 1 by the WDT Control register software reset and cleared by the application program. Note that this bit can only be cleared to 0 by the application program.

Watchdog Timer Operation

The Watchdog Timer operates by providing a device reset when its timer overflows. This means that in the application program and during normal operation the user has to strategically clear the Watchdog Timer before it overflows to prevent the Watchdog Timer from executing a reset. This is done using the clear watchdog instructions. If the program malfunctions for whatever reason, jumps to an unknown location, or enters an endless loop, these clear instructions will not be executed in the correct manner, in which case the Watchdog Timer will overflow and reset the device. With regard to the Watchdog Timer enable/disable function, there are also five bits, WE4~WE0, in the WDTC register to offer additional enable/disable and reset control of the Watchdog Timer. The WDT function will be disabled when the WE4~WE0 bits are set to a value of 10101B. The WDT function will be enabled if the WE4~WE0 bits value is equal to 01010B. If the WE4~WE0 bits are set to any other values by the environmental noise, except 01010B and 10101B, it will reset the device after 2~3 LIRC clock cycles. After power on these bits will have the value of 01010B.

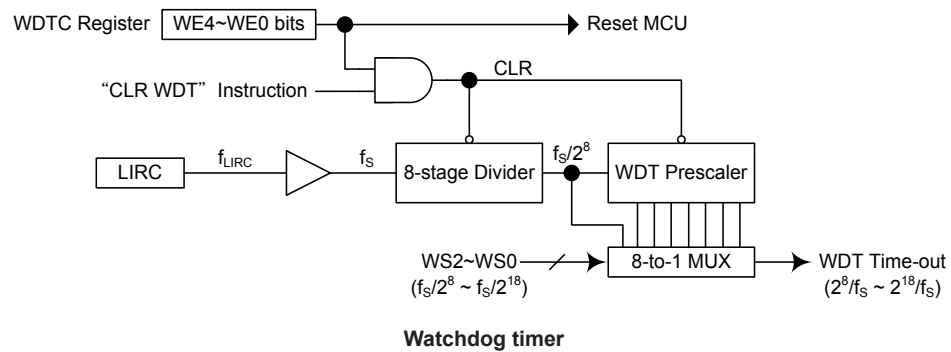
WDT Configuration Option	WE4 ~ WE0 Bits	WDT Function
Application Program Enabled	10101B	Disable
	01010B	Enable
	Any other value	Reset MCU

Watchdog Timer Enable/Disable Control

Under normal program operation, a Watchdog Timer time-out will initialise a device reset and set the status bit TO. However, if the system is in the Power Down Mode, when a Watchdog Timer time-out occurs, the TO bit in the status register will be set and only the Program Counter and Stack Pointer will be reset. Three methods can be adopted to clear the contents of the Watchdog Timer. The first is a WDT reset, which means a certain value except 01010B and 10101B written into the WE4~WE0 bit field, the second is using the Watchdog Timer software clear instructions and the third is via a HALT instruction.

There is only one method of using software instruction to clear the Watchdog Timer. That is to use the single “CLR WDT” instruction to clear the WDT.

The maximum time out period is when the 2^{18} division ratio is selected. As an example, with a 32kHz LIRC oscillator as its source clock, this will give a maximum watchdog period of around 8 second for the 2^{18} division ratio, and a minimum timeout of 7.8ms for the 2^8 division ratio.



Reset and Initialisation

A reset function is a fundamental part of any microcontroller ensuring that the device can be set to some predetermined condition irrespective of outside parameters. The most important reset condition is after power is first applied to the microcontroller. In this case, internal circuitry will ensure that the microcontroller, after a short delay, will be in a well defined state and ready to execute the first program instruction. After this power-on reset, certain important internal registers will be set to defined states before the program commences.

Another type of reset is when the Watchdog Timer overflows and resets the microcontroller. All types of reset operations result in different register conditions being setup. Another reset exists in the form of a Low Voltage Reset, LVR, where a full reset, is implemented in situations where the power supply voltage falls below a certain threshold.

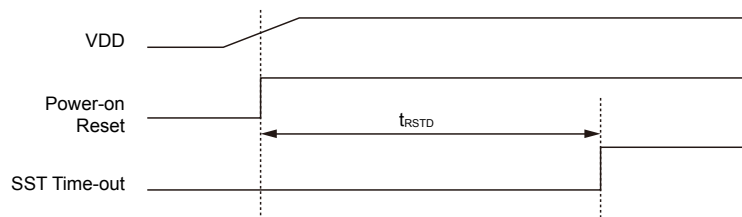
Reset Functions

There are four ways in which a microcontroller reset can occur, through events occurring both internally and externally:

- Power-on Reset

The most fundamental and unavoidable reset is the one that occurs after power is first applied to the microcontroller. As well as ensuring that the Program Memory begins execution from the first memory address, a power-on reset also ensures that certain other registers are preset to known conditions. All the I/O port and port control registers will power up in a high condition ensuring that all pins will be first set to inputs.

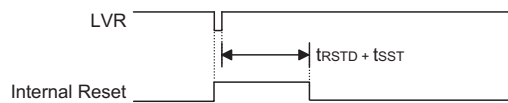
The microcontroller has an internal RC reset function, due to unstable power on conditions. This time delay created by the RC network ensures the state of the POR remains low for an extended period while the power supply stabilizes. During this time, normal operation of the microcontroller is inhibited. After the state of the POR reaches a certain voltage value, the reset delay time t_{POR} is invoked to provide an extra delay time after which the microcontroller can begin normal operation.



Power-On Reset Timing Chart

- Low Voltage Reset – LVR

The microcontroller contains a low voltage reset circuit in order to monitor the supply voltage of the device. The LVR function is always enabled with a specific LVR voltage V_{LVR} . If the supply voltage of the device drops to within a range of $0.9V \sim V_{LVR}$ such as might occur when changing the battery, the LVR will automatically reset the device internally and the LVRF bit in the CTRL1 register will also be set to 1. For a valid LVR signal, a low supply voltage, i.e., a voltage in the range between $0.9V \sim V_{LVR}$ must exist for a time greater than that specified by t_{LVR} in the A.C. characteristics. If the low supply voltage state does not exceed this value, the LVR will ignore the low supply voltage and will not perform a reset function. The actual V_{LVR} value can be selected by the LVS bits in the LVRC register. If the LVS7~LVS0 bits have any other value, which may perhaps occur due to adverse environmental conditions such as noise, the LVR will reset the device after 2~3 LIRC clock cycles. When this happens, the LRF bit in the CTRL1 register will be set to 1. After power on the register will have the value of 01010101B. Note that the LVR function will be automatically disabled when the device enters the power down mode.



Low Voltage Reset Timing Chart

- LVRC Register

Bit	7	6	5	4	3	2	1	0
Name	LVS7	LVS6	LVS5	LVS4	LVS3	LVS2	LVS1	LVS0
R/W	R/W	R/W	R/W	R/W	R	R	R/W	R/W
POR	0	1	0	1	0	1	0	1

Bit 7~0 **LVS7~LVS0:** LVR voltage select
 01010101: 2.1V
 00110011: 2.55V
 10011001: 3.15V
 10101010: 3.55V

Any other values: Generates MCU reset – register is reset to POR value

When an actual low voltage condition occurs, as specified by one of the four defined LVR voltage values above, an MCU reset will be generated. The reset operation will be activated after 2~3 LIRC clock cycles. In this situation the register contents will remain the same after such a reset occurs.

Any register value, other than the four defined register values above, will also result in the generation of an MCU reset. The reset operation will be activated after 2~3 LIRC clock cycles. However in this situation the register contents will be reset to the POR value.

• **CTRL1 Register**

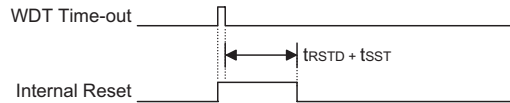
Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	LVRF	LRF	WRF
R/W	—	—	—	—	—	R/W	R/W	R/W
POR	—	—	—	—	—	x	0	0

"x": unknown

- Bit 7~3 Unimplemented, read as "0"
- Bit 2 **LVRF**: LVR function reset flag
 0: Not occurred
 1: Occurred
 This bit is set to 1 when a specific Low Voltage Reset situation condition occurs. This bit can only be cleared to 0 by the application program.
- Bit 1 **LRF**: LVR Control register software reset flag
 0: Not occurred
 1: Occurred
 This bit is set to 1 if the LVRC register contains any non defined LVR voltage register values. This in effect acts like a software reset function. This bit can only be cleared to 0 by the application program.
- Bit 0 **WRF**: WDT Control register software reset flag
 Describe elsewhere.

• **Watchdog Time-out Reset during Normal Operation**

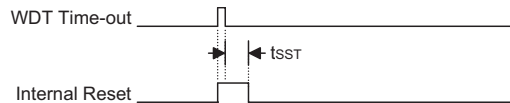
The Watchdog time-out Reset during normal operation is the same as a hardware power-on reset except that the Watchdog time-out flag TO will be set to "1".



WDT Time-out Reset during Normal Operation Timing Chart

• **Watchdog Time-out Reset during Power Down Mode**

The Watchdog time-out Reset during Power Down Mode is a little different from other kinds of reset. Most of the conditions remain unchanged except that the Program Counter and the Stack Pointer will be cleared to "0" and the TO flag will be set to "1". Refer to the A.C. Characteristics for t_{SST} details.



WDT Time-out Reset during Power Down Mode Timing Chart

Note: The t_{SST} is 2 clock cycles for HIRC. The t_{SST} is 1024 clock cycles for HXT.

Reset Initial Conditions

The different types of reset described affect the reset flags in different ways. These flags, known as PDF and TO are located in the status register and are controlled by various microcontroller operations, such as the Power Down Mode function or Watchdog Timer. The reset flags are shown in the table:

TO	PDF	RESET Conditions
0	0	Power-on reset
u	u	LVR reset during NORMAL Mode operation
1	u	WDT time-out reset during NORMAL Mode operation
1	1	WDT time-out reset during Power down Mode operation

“u” unchanged

The following table indicates the way in which the various components of the microcontroller are affected after a power-on reset occurs.

Item	Condition After RESET
Program Counter	Reset to zero
Interrupts	All interrupts will be disabled
Prescaler, Divider	Cleared
WDT	Clear after reset, WDT begins counting
Timer/Event Counter	Timer Counter will be turned off
Input/Output Ports	I/O ports will be setup as inputs
Stack Pointer	Stack Pointer will point to the top of the stack

The different kinds of resets all affect the internal registers of the microcontroller in different ways. To ensure reliable continuation of normal program execution after a reset occurs, it is important to know what condition the microcontroller is in after a particular reset occurs. The following table describes how each type of reset affects the microcontroller internal registers.

Register	HT48R016	HT48R017	Power On Reset	LVR Reset (Normal Operation)	WDT Time-out (Normal Operation)	WDT Time-out (HALT)*
MP0	•	•	x x x x x x x x	x x x x x x x x	x x x x x x x x	u u u u u u u u
MP1	•	•	x x x x x x x x	x x x x x x x x	x x x x x x x x	u u u u u u u u
ACC	•	•	x x x x x x x x	u u u u u u u u	u u u u u u u u	u u u u u u u u
TBLP	•	•	x x x x x x x x	u u u u u u u u	u u u u u u u u	u u u u u u u u
TBLH	•	•	x x x x x x x x	u u u u u u u u	u u u u u u u u	u u u u u u u u
WDTC	•	•	0 1 0 1 0 0 1 1	0 1 0 1 0 0 1 1	0 1 0 1 0 0 1 1	u u u u u u u u
STATUS	•	•	- - 0 0 x x x x	- - u u u u u u	- - 1 u u u u u	- - 1 1 u u u u
INTC	•	•	- 0 0 0 0 0 0 0 0	- 0 0 0 0 0 0 0 0	- 0 0 0 0 0 0 0 0	- u u u u u u u
TMR	•	•	x x x x x x x x	x x x x x x x x	x x x x x x x x	u u u u u u u u
TMRC	•	•	0 0 - 0 1 0 0 0 0	0 0 - 0 1 0 0 0 0	0 0 - 0 1 0 0 0 0	u u - u u u u u
PA	•	•	1 1 1 1 1 1 1 1	1 1 1 1 1 1 1 1	1 1 1 1 1 1 1 1	u u u u u u u u
PAC	•	•	1 1 1 1 1 1 1 1	1 1 1 1 1 1 1 1	1 1 1 1 1 1 1 1	u u u u u u u u
PAPU	•	•	- 0 0 0 0 0 0 0 0	- 0 0 0 0 0 0 0 0	- 0 0 0 0 0 0 0 0	- u u u u u u u
PAWK	•	•	0 0 0 0 0 0 0 0	0 0 0 0 0 0 0 0	0 0 0 0 0 0 0 0	u u u u u u u u
PB	•	•	- - 1 1 1 1 1 1 1 1	- - 1 1 1 1 1 1 1 1	- - 1 1 1 1 1 1 1 1	- - u u u u u u
PBC	•	•	- - 1 1 1 1 1 1 1 1	- - 1 1 1 1 1 1 1 1	- - 1 1 1 1 1 1 1 1	- - u u u u u u
PBPU	•	•	- - 0 0 0 0 0 0 0 0	- - 0 0 0 0 0 0 0 0	- - 0 0 0 0 0 0 0 0	- - u u u u u u
CTRL0	•	•	1 0 0 0 - - - -	1 0 0 0 - - - -	1 0 0 0 - - - -	u u u u - - - -
CTRL1	•	•	- - - - - x 0 0	- - - - - u u u	- - - - - u u u	- - - - - u u u
PFDCTRL	•	•	- - - - - - 0 0	- - - - - - 0 0	- - - - - - 0 0	- - - - - - u u
LVRC	•	•	0 1 0 1 0 1 0 1	0 1 0 1 0 1 0 1	0 1 0 1 0 1 0 1	u u u u u u u u

Note: “-” not implement

“u” means “unchanged”

“x” means “unknown”

Input/Output Ports

Holtek microcontrollers offer considerable flexibility on their I/O ports. Most pins can have either an input or output designation under user program control. Additionally, as there are pull-high resistors and wake-up software configurations, the user is provided with an I/O structure to meet the needs of a wide range of application possibilities.

For input operation, these ports are non-latching, which means the inputs must be ready at the T2 rising edge of instruction “MOV A, [m]”, where m denotes the port address. For output operation, all the data is latched and remains unchanged until the output latch is rewritten.

Pull-high Resistors

Many product applications require pull-high resistors for their switch inputs usually requiring the use of an external resistor. To eliminate the need for these external resistors, all I/O pins, when configured as an input have the capability of being connected to an internal pull-high resistor. These pull-high resistors are selected using registers PAPU, PBPU located in the Data Memory. The pull-high resistors are implemented using weak PMOS transistors. Note that pin PA7 does not have a pull-high resistor selection.

Port A Wake-up

If the HALT instruction is executed, the device will enter the Power Down Mode, where the system clock will stop resulting in power being conserved, a feature that is important for battery and other low-power applications. Various methods exist to wake-up the microcontroller, one of which is to change the logic condition on one of the PA0~PA7 pins from high to low. After a HALT instruction forces the microcontroller into entering the Idle/Sleep Mode, the processor will remain idle or in a low-power state until the logic condition of the selected wake-up pin on Port A changes from high to low. This function is especially suitable for applications that can be woken up via external switches. Note that pins PA0 to PA7 can be selected individually to have this wake-up feature using an internal register known as PAWK, located in the Data Memory.

PAWK, PAC, PAPU, PBC, PBPU Register

Register Name	Bit							
	7	6	5	4	3	2	1	0
PAWK	PAWK7	PAWK6	PAWK5	PAWK4	PAWK3	PAWK2	PAWK1	PAWK0
PAC	PAC7	PAC6	PAC5	PAC4	PAC3	PAC2	PAC1	PAC0
PAPU	—	PAPU6	PAPU5	PAPU4	PAPU3	PAPU2	PAPU1	PAPU0
PBC	—	—	PBC5	PBC4	PBC3	PBC2	PBC1	PBC0
PBPU	—	—	PBPU5	PBPU4	PBPU3	PBPU2	PBPU1	PBPU0

“—” Unimplemented, read as “0”

PAWKn: PA wake-up function enable

0: disable

1: enable

PACn/PBCn: I/O type selection

0: output

1: input

PAPUn/PBPU: Pull-high function enable

0: disable

1: enable

I/O Port Control Registers

Each Port has its own control register, known as PAC, PBC which controls the input/output configuration. With this control register, each I/O pin with or without pull-high resistors can be reconfigured dynamically under software control. For the I/O pin to function as an input, the corresponding bit of the control register must be written as a “1”. This will then allow the logic state of the input pin to be directly read by instructions. When the corresponding bit of the control register is written as a “0”, the I/O pin will be setup as a CMOS output. If the pin is currently setup as an output, instructions can still be used to read the output register. However, it should be noted that the program will in fact only read the status of the output data latch and not the actual logic status of the output pin.

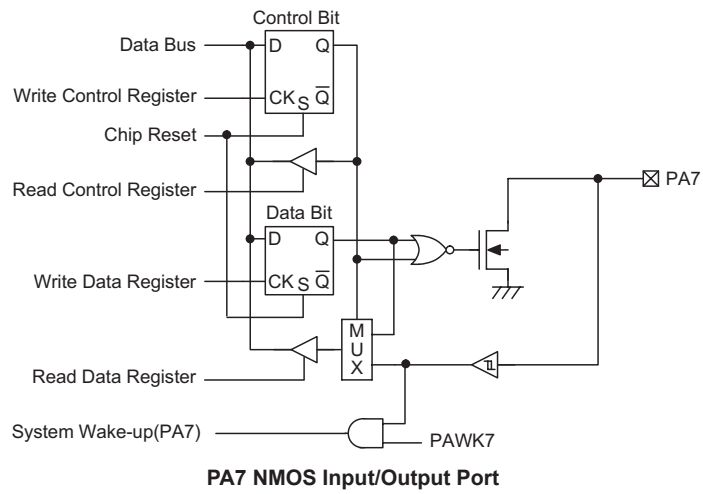
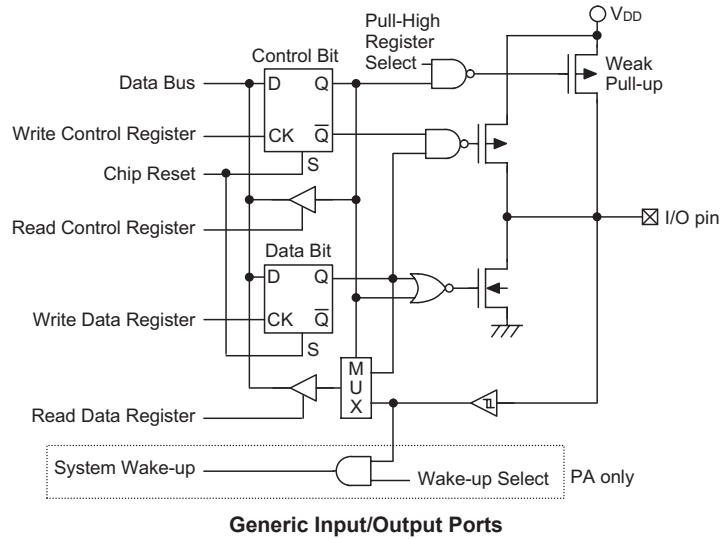
Pin-shared Functions

The flexibility of the microcontroller range is greatly enhanced by the use of pins that have more than one function. Limited numbers of pins can force serious design constraints on designers but by supplying pins with multi-functions, many of these difficulties can be overcome. For some pins, the chosen function of the multi-function I/O pins is set by configuration options while for others the function is set by application program control.

- **External Interrupt Input**
The external interrupt pin, INT, is pin-shared with an I/O pin. To use the pin as an external interrupt input the correct bits in the INTC0 register must be programmed. The pin must also be setup as an input by setting the PAC3 bit in the Port Control Register. A pull-high resistor can also be selected via the appropriate port pull-high resistor register. Note that even if the pin is setup as an external interrupt input the I/O function still remains.
- **External Timer/Event Counter Input**
The Timer/Event Counter pins, TC are pin-shared with I/O pin. For the shared pin to be used as Timer/Event Counter input, the Timer/Event Counter must be configured to be in the Event Counter or Pulse Width Capture Mode. This is achieved by setting the appropriate bits in the Timer/Event Counter Control Register. The pin must also be setup as input by setting the appropriate bit in the Port Control Register. Pull-high resistor options can also be selected using the port pull-high resistor registers. Note that even if the pin is setup as an external timer input the I/O function still remains.
- **PFD Output**
The PFD complementary output pair is pin-shared with I/O pins. The output function of these pins is chosen using the PFDCTRL register. Note that the corresponding bit of the port control register, must setup the pin as an output to enable the PFD output. If the port control register has setup the pin as an input, then the pin will function as a normal logic input with the usual pull-high selection, even if the PFD function has been selected.

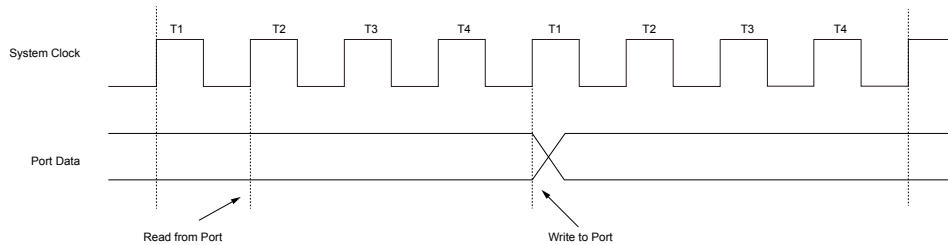
I/O Pin Structures

The diagrams illustrate the I/O pin internal structures. As the exact logical construction of the I/O pin may differ from these drawings, they are supplied as a guide only to assist with the functional understanding of the I/O pins.



Programming Considerations

Within the user program, one of the things first to consider is port initialisation. After a reset, all of the I/O data and port control registers will be set to high. This means that all I/O pins will be defaulted to an input state, the level of which depends on the other connected circuitry and whether pull-high selections have been chosen. If the port control registers are then programmed to setup some pins as outputs, these output pins will have an initial high output value unless the associated port data registers are first programmed. Selecting which pins are inputs and which are outputs can be achieved byte-wide by loading the correct values into the appropriate port control register or by programming individual bits in the port control register using the “SET [m].i” and “CLR [m].i” instructions. Note that when using these bit control instructions, a read-modify-write operation takes place. The microcontroller must first read in the data on the entire port, modify it to the required new bit values and then rewrite this data back to the output ports.



Read Modify Write Timing

Pins PA0 to PA7 each have a wake-up functions, selected via the PAWK register. When the device is in the Idle/Sleep Mode, various methods are available to wake the device up. One of these is a high to low transition of any of the these pins. Single or multiple pins on Port A can be setup to have this function.

Timer/Event Counter

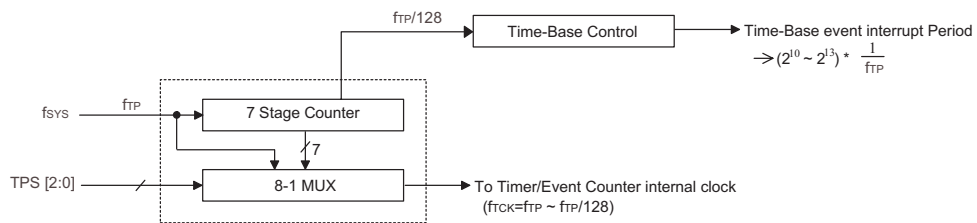
The provision of timers form an important part of any microcontroller, giving the designer a means of carrying out time related functions. The devices contain from one to two count-up timers of 8-bit capacity. As the timers have three different operating modes, they can be configured to operate as a general timer, an external event counter or as a pulse width capture device. The provision of an internal prescaler to the clock circuitry on giving added range to the timers.

There are two types of registers related to the Timer/Event Counters. The first is the register that contains the actual value of the timer and into which an initial value can be preloaded. Reading from this register retrieves the contents of the Timer/Event Counter. The second type of associated register is the Timer Control Register which defines the timer options and determines how the timer is to be used. The device can have the timer clock configured to come from the internal clock source. In addition, the timer clock source can also be configured to come from an external timer pin.

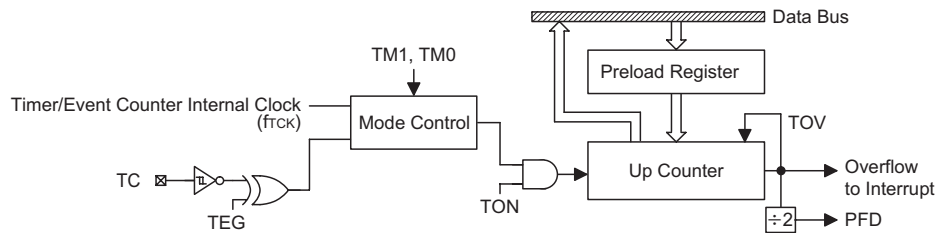
Configuring the Timer/Event Counter Input Clock Source

The Timer/Event Counter clock source can originate from an internal clock or an external pin. The internal clock source is used when the timer is in the timer mode or in the pulse width capture mode. The internal clock source is derived from the system clock, f_{SYS} , and is first divided by a prescaler, the division ratio of which is conditioned by the Timer Control Register bits TPSC0~TPSC2.

An external clock source is used when the timer is in the event counting mode, the clock source being provided on an external timer pin TC. Depending upon the condition of the TEG bit, each high to low, or low to high transition on the external timer pin will increment the counter by one.



Clock Structure



8-bit Timer/Event Counter Structure

Timer Register – TMR

The timer register is a special function register located in the Special Purpose Data Memory and is the place where the actual timer value is stored. The register is known as TMR. The value in the timer register increases by one each time an internal clock pulse is received or an external transition occurs on the external timer pin. The timer will count from the initial value loaded by the preload register to the full count of FFH at which point the timer overflows and an internal interrupt signal is generated. Then the timer value will be reset with the initial preload register value and continue counting.

Note that to achieve a maximum full range count of FFH, all the preload registers must first be cleared to zero. It should be noted that after power-on, the preload registers will be in an unknown condition. Note that if the Timer/Event Counter is in an OFF condition and data is written to its preload register, this data will be immediately written into the actual counter. However, if the counter is enabled and counting, any new data written into the preload data register during this period will remain in the preload register and will only be written into the actual counter the next time an overflow occurs.

Timer Control Register – TMRC

The flexible features of the Holtek microcontroller Timer/Event Counters enable them to operate in three different modes, the options of which are determined by the contents of their respective control register.

The Timer Control Register is known as TMRC. It is the Timer Control Register together with its corresponding timer registers that control the full operation of the Timer/Event Counter. Before the timer can be used, it is essential that the Timer Control Register is fully programmed with the right data to ensure its correct operation, a process that is normally carried out during program initialisation.

To choose which of the three modes the timer is to operate in, either in the timer mode, the event counting mode or the pulse width capture mode, bits 7 and 6 of the Timer Control Register, which are known as the bit pair TM1/TM0, must be set to the required logic levels. The timer-on bit, which is bit 4 of the Timer Control Register and known as TON, provides the basic on/off control of the respective timer. Setting the bit high allows the counter to run. Clearing the bit stops the counter. Bits 0~2 of the Timer Control Register determines the division ratio of the input clock prescaler. The prescaler bit settings have no effect if an external clock source is used. If the timer is in the event count or pulse width capture mode, the active transition edge level type is selected by the logic level of bit 3 of the Timer Control Register which is known as TEG.

TMRC Register

Bit	7	6	5	4	3	2	1	0
Name	TM1	TM0	—	TON	TEG	TPSC2	TPSC1	TPSC0
R/W	R/W	R/W	—	R/W	R/W	R/W	R/W	R/W
POR	0	0	—	0	1	0	0	0

- Bit 7, 6 **TM1, TM0:** Timer/Event Counter operation mode selection
 00: no mode available
 01: event counter mode
 10: timer mode
 11: pulse width capture mode
- Bit 5 Unimplemented, read as “0”
- Bit 4 **TON:** Timer/event counter counting enable
 0: disable
 1: enable
- Bit 3 **TEG:**
 Event counter active edge selection
 1: count on falling edge
 0: count on rising edge
 Pulse Width Capture active edge selection
 1: start counting on rising edge, stop on falling edge
 0: start counting on falling edge, stop on rising edge
- Bit 2~0 **TPSC2, TPSC1, TPSC0:** Timer prescaler rate selection Timer internal clock
 000: f_{TP}
 001: $f_{TP}/2$
 010: $f_{TP}/4$
 011: $f_{TP}/8$
 100: $f_{TP}/16$
 101: $f_{TP}/32$
 110: $f_{TP}/64$
 111: $f_{TP}/128$

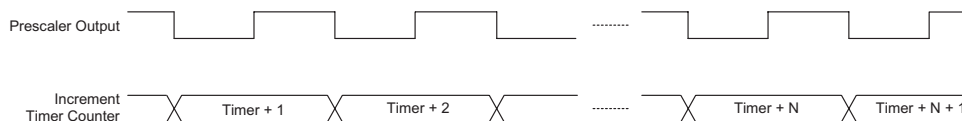
Timer Mode

In this mode, the Timer/Event Counter can be utilised to measure fixed time intervals, providing an internal interrupt signal each time the Timer/Event Counter overflows. To operate in this mode, the Operating Mode Select bit pair, TM1/TM0, in the Timer Control Register must be set to the correct value as shown.

Bit7	Bit6
1	0

Control Register Operating Mode Select Bits for the Timer Mode

In this mode the internal clock is used as the timer clock. The timer input clock source is f_{sys} . However, this timer clock source is further divided by a prescaler, the value of which is determined by the bits TPSC2~TPSC0 in the Timer Control Register. The timer-on bit, TON must be enable the timer to run. Each time an internal clock high to low transition occurs, the timer increments by one. When the timer is full and overflows, an interrupt signal is generated and the timer will reload the value already loaded into the preload register and continue counting. A timer overflow condition and corresponding internal interrupts are two of the wake-up sources. However, the internal interrupt can be disabled by ensuring that the TE bit of the INTC register is reset to zero.



Timer Mode Timing Chart

Event Counter Mode

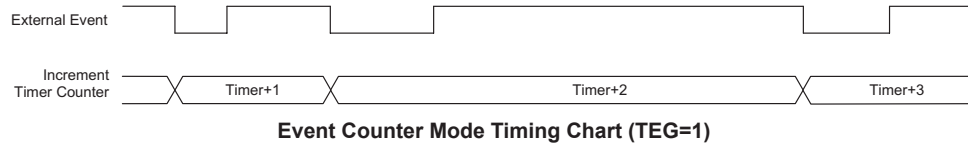
In this mode, a number of externally changing logic events, occurring on the external timer TC pin, can be recorded by the Timer/Event Counter. To operate in this mode, the Operating Mode Select bit pair, TM1/TM0, in the Timer Control Register must be set to the correct value as shown.

Bit7	Bit6
0	1

Control Register Operating Mode Select Bits for the Timer Mode

In this mode, the external timer input, TC, is used as the Timer/Event Counter clock source, however it is not divided by the internal prescaler. After the other bits in the Timer Control Register have been setup, the enable bit TON, which is bit 4 of the Timer Control Register, can be set high to enable the Timer/Event Counter to run. If the Active Edge Select bit, TEG, which is bit 3 of the Timer Control Register, is low, the Timer/Event Counter will increment each time the external timer pin receives a low to high transition. If the TEG is high, the counter will increment each time the external timer pin receives a high to low transition. When it is full and overflows, an interrupt signal is generated and the Timer/Event Counter will reload the value already loaded into the preload register and continue counting. The interrupt can be disabled by ensuring that the Timer/Event Counter Interrupt Enable bit in the corresponding Interrupt Control Register is reset to zero.

As the external timer pin is shared with an I/O pin, to ensure that the pin is configured to operate as an event counter input pin, two things have to happen. The first is to ensure that the Operating Mode Select bits in the Timer Control Register place the Timer/Event Counter in the Event Counting Mode. The second is to ensure that the port control register configures the pin as an input. It should be noted that in the event counting mode, even if the microcontroller is in the Power Down Mode, the Timer/Event Counter will continue to record externally changing logic events on the timer input TC pin. As a result when the timer overflows it will generate a timer interrupt and corresponding wake-up source.



Pulse Width Capture Mode

In this mode, the Timer/Event Counter can be utilised to measure the width of external pulses applied to the external timer pin. To operate in this mode, the Operating Mode Select bit pair, TM1/TM0, in the Timer Control Register must be set to the correct value as shown.

Bit7	Bit6
1	1

Control Register Operating Mode Select Bits for the Pulse Width Capture Mode

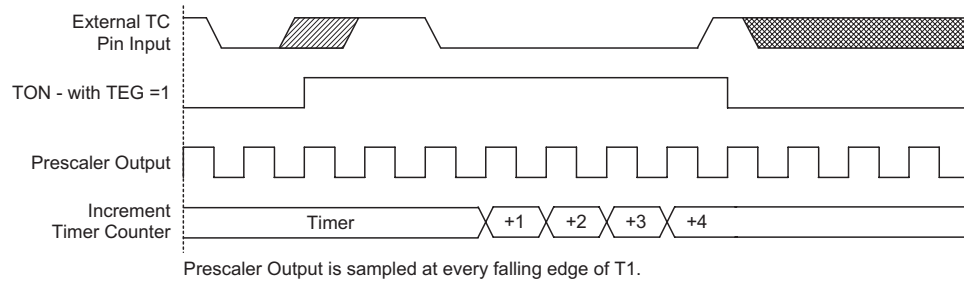
In this mode the internal clock, f_{SYS} , is used as the internal clock for the 8-bit Timer/Event Counter. However, the clock source, f_{SYS} , for the 8-bit timer is further divided by a prescaler, the value of which is determined by the Prescaler Rate Select bits TPSC2~TPSC0, which TON, which is bit 2~0 of the Timer Control Register, can be set high to enable the Timer/Event Counter, however it will not actually start counting until an active edge is received on the external timer pin.

If the Active Edge Select bit TEG, which is bit 3 of the Timer Control Register, is low, once a high to low transition has been received on the external timer pin, the Timer/Event Counter will start counting until the external timer pin returns to its original high level. At this point the enable bit will be automatically reset to zero and the Timer/Event Counter will stop counting. If the Active Edge Select bit is high, the Timer/Event Counter will begin counting once a low to high transition has been received on the external timer pin and stop counting when the external timer pin returns to its original low level. As before, the enable bit will be automatically reset to zero and the Timer/Event Counter will stop counting. It is important to note that in the pulse width capture mode, the enable bit is automatically reset to zero when the external control signal on the external timer pin returns to its original level, whereas in the other two modes the enable bit can only be reset to zero under program control.

The residual value in the Timer/Event Counter, which can now be read by the program, therefore represents the length of the pulse received on the TC pin. As the enable bit has now been reset, any further transitions on the external timer pin will be ignored. The timer cannot begin further pulse width capture until the enable bit is set high again by the program. In this way, single shot pulse measurements can be easily made.

It should be noted that in this mode the Timer/Event Counter is controlled by logical transitions on the external timer pin and not by the logic level. When the Timer/Event Counter is full and overflows, an interrupt signal is generated and the Timer/Event Counter will reload the value already loaded into the preload register and continue counting. The interrupt can be disabled by ensuring that the Timer/Event Counter Interrupt Enable bit in the corresponding Interrupt Control Register, it is reset to zero.

As the TC pin is shared with an I/O pin, to ensure that the pin is configured to operate as a pulse width capture pin, two things have to be implemented. The first is to ensure that the Operating Mode Select bits in the Timer Control Register place the Timer/Event Counter in the pulse width capture mode, the second is to ensure that the port control register configure the pin as an input.



Pulse Width Capture Mode Timing Chart (TEG=1)

Prescaler

Bits TPSC0~TPSC2 of the TMRC register can be used to define a division ratio for the internal clock source of the Timer/Event Counter enabling longer time-out periods to be setup.

PFD Function

The Programmable Frequency Divider provides a means of producing a variable frequency output suitable for applications, such as piezo-buzzer driving or other interfaces requiring a precise frequency generator.

The Timer/Event Counter overflow signal is the clock source for the PFD function, which is controlled by PFDCS bit in PFDCTRL. The PFD clock source comes from the Timer/Event Counter. The output frequency is controlled by loading the required values into the timer prescaler and timer registers to give the required division ratio. The counter will begin to count-up from this preload register value until full, at which point an overflow signal is generated, causing both the PFD outputs to change state. Then the counter will be automatically reloaded with the preload register value and continue counting-up.

If the PFDCTRL register has selected the PFD function, then for PFD output to operate, it is essential for the Port A control register PAC to setup the PFD pins as outputs. The PA0 data bit must be set high to activate the PFD function if only the PFD function is enabled. However, if the PFD and $\overline{\text{PFD}}$ functions both are enabled, the PA0 and PA1 data bits must be set high to activate the PFD and $\overline{\text{PFD}}$ complementary outputs. The output data bits can be used as the on/off control bit for the PFD outputs. Note that the PFD outputs will all be low if the output data bit is cleared to zero.

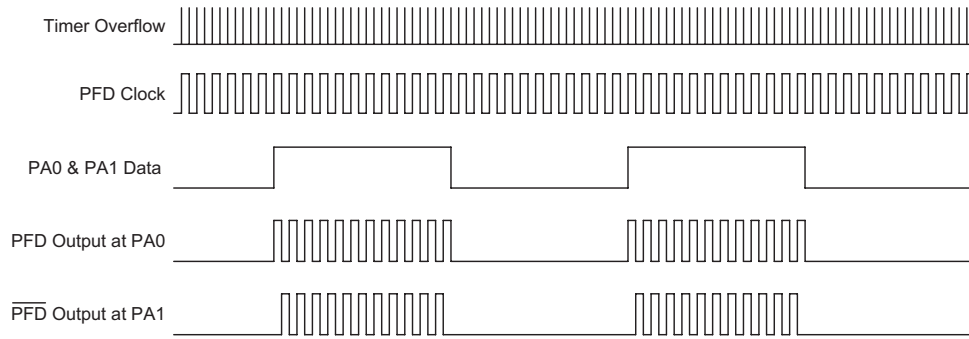
Using this method of frequency generation, and if a crystal oscillator is used for the system clock, very precise values of frequency can be generated.

PFDCTRL Register

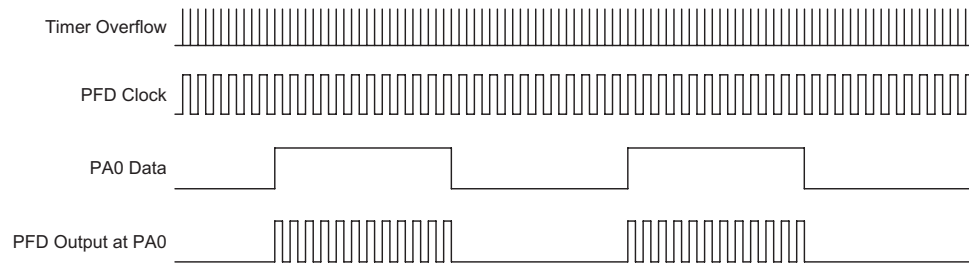
Bit	7	6	5	4	3	2	1	0
Name	—	—	—	—	—	—	PFDEN1	PFDEN0
R/W	—	—	—	—	—	—	R/W	R/W
POR	—	—	—	—	—	—	0	0

Bit7~2 Unimplemented, read as "0"

Bit1~0 **PFDEN1~PFDEN0**: PFD/ $\overline{\text{PFD}}$ enable control
 00: both disabled
 01: Reserved
 10: PFD enabled
 11: PFD and $\overline{\text{PFD}}$ both enabled



PFD Function – Complementary Outputs



PFD Function – Single Output

I/O Interfacing

The Timer/Event Counter, when configured to run in the event counter or pulse width capture mode, requires the use of an external timer pin for its operation. As this pin is a shared pin it must be configured correctly to ensure that it is setup for use as a Timer/Event Counter input pin. This is achieved by ensuring that the mode selects bits in the Timer/Event Counter control register, either the event counter or pulse width capture mode. Additionally the corresponding Port Control Register bit must be set high to ensure that the pin is setup as an input. Any pull-high resistor connected to this pin will remain valid even if the pin is used as a Timer/Event Counter input.

Programming Considerations

When configured to run in the timer mode, the internal system clock is used as the timer clock source and is therefore synchronised with the overall operation of the microcontroller. In this mode when the appropriate timer register is full, the microcontroller will generate an internal interrupt signal directing the program flow to the respective internal interrupt vector. For the pulse width capture mode, the internal system clock is also used as the timer clock source but the timer will only run when the correct logic condition appears on the external timer input pin. As this is an external event and not synchronised with the internal timer clock, the microcontroller will only see this external event when the next timer clock pulse arrives. As a result, there may be small differences in measured values requiring programmers to take this into account during programming. The same applies if the timer is configured to be in the event counting mode, which again is an external event and not synchronised with the internal system or timer clock.

When the Timer/Event Counter is read, or if data is written to the preload register, the clock is inhibited to avoid errors, however as this may result in a counting error, this should be taken into account by the programmer. Care must be taken to ensure that the timers are properly initialised before using them for the first time. The associated timer enable bits in the interrupt control register must be properly set otherwise the internal interrupt associated with the timer will remain inactive. The edge select, timer mode and clock source control bits in timer control register must also be correctly set to ensure the timer is properly configured for the required application. It is also important to ensure that an initial value is first loaded into the timer registers before the timer is switched on; this is because after power-on the initial values of the timer registers are unknown. After the timer has been initialized the timer can be turned on and off by controlling the enable bit in the timer control register.

When the Timer/Event Counter overflows, its corresponding interrupt request flag in the interrupt control register will be set. If the Timer/Event Counter interrupt is enabled this will in turn generate an interrupt signal. However irrespective of whether the interrupts are enabled or not, a Timer/Event Counter overflow will also generate a wake-up signal if the device is in a Power-down condition. This situation may occur if the Timer/Event Counter is in the Event Counting Mode and if the external signal continues to change state. In such a case, the Timer/Event Counter will continue to count these external events and if an overflow occurs the device will be woken up from its Power-down condition. To prevent such a wake-up from occurring, the timer interrupt request flag should first be set high before issuing the "HALT" instruction to enter the Power Down Mode.

Timer Program Example

The program shows how the Timer/Event Counter registers are setup along with how the interrupts are enabled and managed. Note how the Timer/Event Counter is turned on, by setting bit 4 of the Timer Control Register. The Timer/Event Counter can be turned off in a similar way by clearing the same bit. This example program sets the Timer/Event Counters to be in the timer mode, which uses the internal system clock as their clock source.

Timer Programming Example

```
org 04h          ; external interrupt vector
org 08h          ; Timer Counter interrupt vector
jmp tmrint      ; jump here when Timer overflows
:
:
org 20h          ; main program
:
:
:                ; internal Timer interrupt routine
tmrint:
:
:                ; Timer main program placed here
:
:
begin:
:                ; setup Timer registers
mov a,09bh      ; setup Timer preload value
mov tmr,a
mov a,081h      ; setup Timer control register
mov tmrc,a      ; timer mode and prescaler set to /2
:                ; setup interrupt register
mov a,00dh      ; enable master interrupt and both timer interrupts
mov intc,a
:
:
set tmrc.4      ; start Timer
:
:
```

Time Base

The device includes a Time Base function which is used to generate a regular time interval signal. The Time Base time interval magnitude is determined using an internal 13 stage counter sets the division ratio of the clock source. This division ratio is controlled by both the TBSEL0 and TBSEL1 bits in the CTRL0 register. The clock source is derived from the system clock f_{SYS} .

When the Time Base time-out condition occurs, a Time Base interrupt signal will be generated. It should be noted that as the Time Base clock source is the same as the Timer/Event Counter clock source, care should be taken when programming.

Interrupts

Interrupts are an important part of any microcontroller system. When an external event or an internal function such as a Timer/Event Counter or Time Base require microcontroller attention, their corresponding interrupt will enforce a temporary suspension of the main program allowing the microcontroller to direct attention to their respective needs.

The devices contain one external interrupt and two internal interrupts. The external interrupt is controlled by the action of the external interrupt pin, while the internal interrupt is controlled by the Timer/Event Counters and Time Base function.

Interrupt Register

Overall interrupt control, which means interrupt enabling and request flag setting, is controlled by using the register, INTC. By controlling the appropriate enable bits in the register each individual interrupt can be enabled or disabled. Also when an interrupt occurs, the corresponding request flag will be set by the microcontroller. The global enable flag cleared to zero will disable all interrupts.

INTC Register

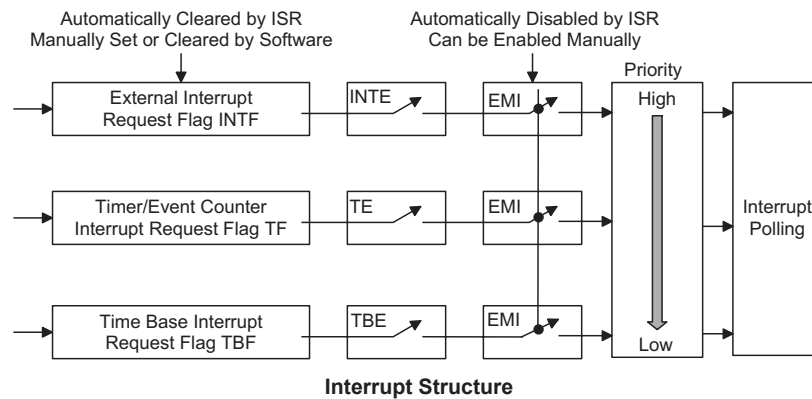
Bit	7	6	5	4	3	2	1	0
Name	—	TBF	TF	INTF	TBE	TE	INTE	EMI
R/W	—	R/W	R/W	R/W	R/W	R/W	R/W	R/W
POR	—	0	0	0	0	0	0	0

- Bit 7 Unimplemented, read as "0"
- Bit 6 **TBF:** Time Base event interrupt request flag
0: inactive
1: active
- Bit 5 **TF:** Timer/Event Counter interrupt request flag
0: inactive
1: active
- Bit 4 **INTF:** External interrupt request flag
0: inactive
1: active
- Bit 3 **TBE:** Time Base event interrupt enable
0: disable
1: enable
- Bit 2 **TE:** Timer/Event Counter interrupt enable
0: disable
1: enable
- Bit 1 **INTE:** External interrupt enable
0: disable
1: enable
- Bit 0 **EMI:** Master interrupt global enable
0: disable
1: enable

Interrupt Operation

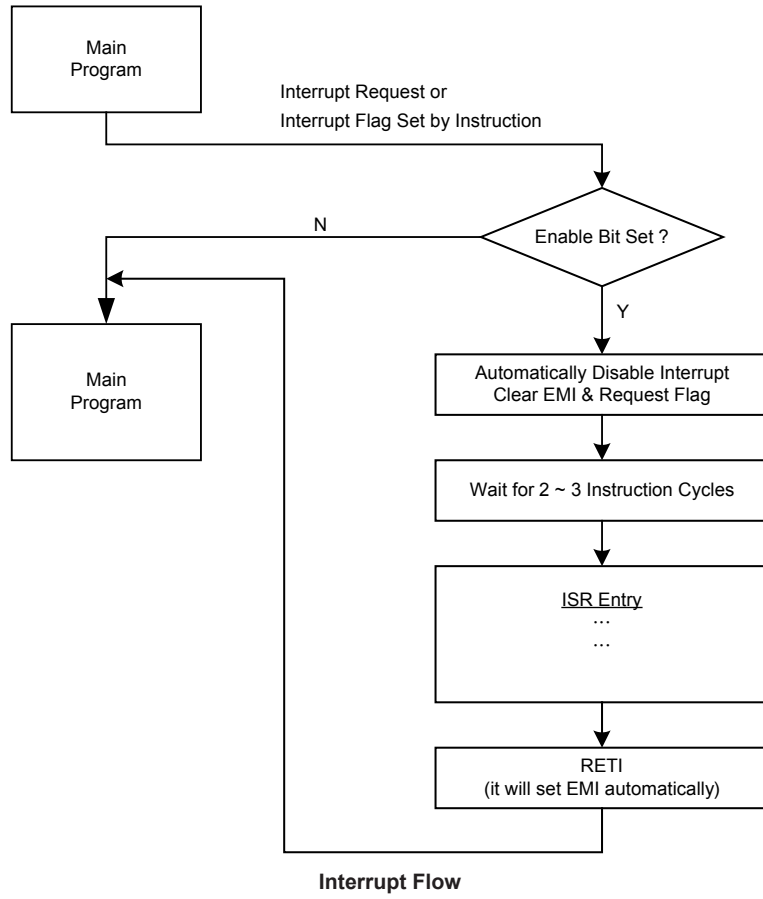
A Timer/Event Counter overflow, a Time Base event, or an active edge on the external interrupt pin will all generate an interrupt request by setting their corresponding request flag, if their appropriate interrupt enable bit is set. When this happens, the Program Counter, which stores the address of the next instruction to be executed, will be transferred onto the stack. The Program Counter will then be loaded with a new address which will be the value of the corresponding interrupt vector. The microcontroller will then fetch its next instruction from this interrupt vector. The instruction at this vector will usually be a JMP statement which will jump to another section of program which is known as the interrupt service routine. Here is located the code to control the appropriate interrupt. The interrupt service routine must be terminated with a RETI instruction, which retrieves the original Program Counter address from the stack and allows the microcontroller to continue with normal execution at the point where the interrupt occurred.

The various interrupt enable bits, together with their associated request flags, are shown in the following diagram with their order of priority.



Once an interrupt subroutine is serviced, all the other interrupts will be blocked, as the EMI bit will be cleared automatically. This will prevent any further interrupt nesting from occurring. However, if other interrupt requests occur during this interval, although the interrupt will not be immediately serviced, the request flag will still be recorded. If an interrupt requires immediate servicing while the program is already in another interrupt service routine, the EMI bit should be set after entering the routine, to allow interrupt nesting. If the stack is full, the interrupt request will not be acknowledged, even if the related interrupt is enabled, until the Stack Pointer is decremented. If immediate service is desired, the stack must be prevented from becoming full.

When an interrupt request is generated it takes 2 or 3 instruction cycles before the program jumps to the interrupt vector. If the device is in the Sleep or Idle Mode and is woken up by an interrupt request then it will take 3 cycles before the program jumps to the interrupt vector.



Interrupt Priority

Interrupts, occurring in the interval between the rising edges of two consecutive T2 pulses, will be serviced on the latter of the two T2 pulses, if the corresponding interrupts are enabled. In case of simultaneous requests, the following table shows the priority that is applied. These can be masked by resetting the EMI bit.

HT48R016/HT48R017		
Interrupt Source	Priority	Vector
External interrupt	1	04H
Timer/Event Counter overflow	2	08H
Time Base overflow	3	0CH

In cases where both external and internal interrupts are enabled and where an external and internal interrupt occurs simultaneously, the external interrupt will always have priority and will therefore be serviced first. Suitable masking of the individual interrupts using the interrupt registers can prevent simultaneous occurrences.

External Interrupt

For an external interrupt to occur, the global interrupt enable bit, EMI, and external interrupt enable bit, INTE, must first be set. An actual external interrupt will take place when the external interrupt request flag, INTF, is set, a situation that will occur when an edge transition appears on the external INT line. The type of transition that will trigger an external interrupt, whether high to low, low to high or both is determined by the INTEG0 and INTEG1 bits, which are bits 6 and 7 respectively, in the CTRL0 control register. These two bits can also disable the external interrupt function.

INTEG1	INTEG0	Edge Trigger Type
0	0	External interrupt disable
0	1	Rising edge Trigger
1	0	Falling edge Trigger
1	1	Both edge Trigger

The external interrupt pin is pin-shared with the I/O pin PA3 and can only be configured as an external interrupt pin if the corresponding external interrupt enable bit in the INTC0 register has been set and the edge trigger type has been selected using the CTRL0 register. The pin must also be setup as an input by setting the corresponding PAC.3 bit in the port control register. When the interrupt is enabled, the stack is not full and a transition appears on the external interrupt pin, a subroutine call to the external interrupt vector at location 04H, will take place. When the interrupt is serviced, the external interrupt request flag, INTF, will be automatically reset and the EMI bit will be automatically cleared to disable other interrupts. Note that any pull-high resistor connections on this pin will remain valid even if the pin is used as an external interrupt input.

Timer/Event Counter Interrupt

For a Timer/Event Counter interrupt to occur, the global interrupt enable bit, EMI, and the corresponding timer interrupt enable bit, TE, must first be set. An actual Timer/Event Counter interrupt will take place when the Timer/Event Counter request flag, TF, is set, a situation that will occur when the relevant Timer/Event Counter overflows. When the interrupt is enabled, the stack is not full and a Timer/Event Counter overflow occurs, a subroutine call to the relevant timer interrupt vector, will take place. When the interrupt is serviced, the timer interrupt request flag, TF, will be automatically reset and the EMI bit will be automatically cleared to disable other interrupts.

Time Base Interrupts

For a time base interrupt to occur the global interrupt enable bit EMI and the corresponding interrupt enable bit TBE, must first be set. An actual Time Base interrupt will take place when the time base request flag TBF is set, a situation that will occur when the Time Base overflows. When the interrupt is enabled, the stack is not full and a time base overflow occurs a subroutine call to time base vector will take place. When the interrupt is serviced, the time base interrupt flag. TBF will be automatically reset and the EMI bit will be automatically cleared to disable other interrupts.

Interrupt Wake-up Function

Each of the interrupt functions has the capability of waking up the microcontroller when in the Power Down Mode. A wake-up is generated when an interrupt request flag changes from low to high and is independent of whether the interrupt is enabled or not. Therefore, even though the device is in the Power Down Mode and its system oscillator stopped, situations such as external edge transitions on the external interrupt pins, a low power supply voltage or may cause their respective interrupt flag to be set high and consequently generate an interrupt. Care must therefore be taken if spurious wake-up situations are to be avoided. If an interrupt wake-up function is to be disabled then the corresponding interrupt request flag should be set high before the device enters the Power Down Mode. The interrupt enable bits have no effect on the interrupt wake-up function.

Programming Considerations

By disabling the relevant interrupt enable bits, a requested interrupt can be prevented from being serviced, however, once an interrupt request flag is set, it will remain in this condition in the interrupt register until the corresponding interrupt is serviced or until the request flag is cleared by a software instruction.

It is recommended that programs do not use the “CALL” instruction within the interrupt service subroutine. Interrupts often occur in an unpredictable manner or need to be serviced immediately. If only one stack is left and the interrupt is not well controlled, the original control sequence will be damaged once a "CALL subroutine" is executed in the interrupt subroutine.

All of these interrupts have the capability of waking up the processor when in the Power Down Mode.

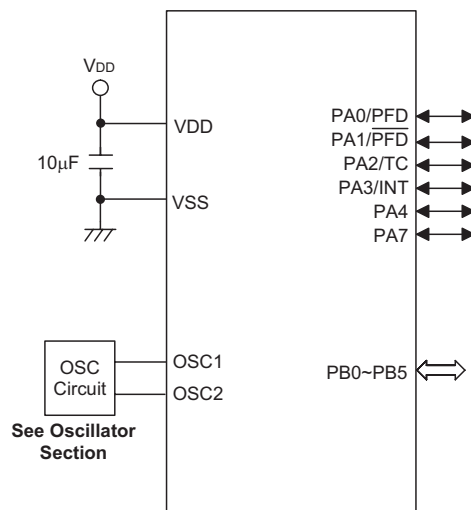
Only the Program Counter is pushed onto the stack. If the contents of the register or status register are altered by the interrupt service program, which may corrupt the desired control sequence, then the contents should be saved in advance.

Configuration Options

Configuration options refer to certain options within the MCU that are programmed into the device during the programming process. During the development process, these options are selected using the HT-IDE software development tools. As these options are programmed into the device using the hardware programming tools, once they are selected they cannot be changed later using the application program. All options must be defined for proper system function, the details of which are shown in the table.

No.	Options
1	System oscillator configuration: HXT or HIRC.
2	HIRC Frequency Selection: 4MHz, 8MHz or 12MHz.

Application Circuits



Instruction Set

Introduction

Central to the successful operation of any microcontroller is its instruction set, which is a set of program instruction codes that directs the microcontroller to perform certain operations. In the case of Holtek microcontrollers, a comprehensive and flexible set of over 60 instructions is provided to enable programmers to implement their application with the minimum of programming overheads.

For easier understanding of the various instruction codes, they have been subdivided into several functional groupings.

Instruction Timing

Most instructions are implemented within one instruction cycle. The exceptions to this are branch, call, or table read instructions where two instruction cycles are required. One instruction cycle is equal to 4 system clock cycles, therefore in the case of an 8MHz system oscillator, most instructions would be implemented within 0.5 μ s and branch or call instructions would be implemented within 1 μ s. Although instructions which require one more cycle to implement are generally limited to the JMP, CALL, RET, RETI and table read instructions, it is important to realize that any other instructions which involve manipulation of the Program Counter Low register or PCL will also take one more cycle to implement. As instructions which change the contents of the PCL will imply a direct jump to that new address, one more cycle will be required. Examples of such instructions would be "CLR PCL" or "MOV PCL, A". For the case of skip instructions, it must be noted that if the result of the comparison involves a skip operation then this will also take one more cycle, if no skip is involved then only one cycle is required.

Moving and Transferring Data

The transfer of data within the microcontroller program is one of the most frequently used operations. Making use of three kinds of MOV instructions, data can be transferred from registers to the Accumulator and vice-versa as well as being able to move specific immediate data directly into the Accumulator. One of the most important data transfer applications is to receive data from the input ports and transfer data to the output ports.

Arithmetic Operations

The ability to perform certain arithmetic operations and data manipulation is a necessary feature of most microcontroller applications. Within the Holtek microcontroller instruction set are a range of add and subtract instruction mnemonics to enable the necessary arithmetic to be carried out. Care must be taken to ensure correct handling of carry and borrow data when results exceed 255 for addition and less than 0 for subtraction. The increment and decrement instructions INC, INCA, DEC and DECA provide a simple means of increasing or decreasing by a value of one of the values in the destination specified.

Logical and Rotate Operations

The standard logical operations such as AND, OR, XOR and CPL all have their own instruction within the Holtek microcontroller instruction set. As with the case of most instructions involving data manipulation, data must pass through the Accumulator which may involve additional programming steps. In all logical data operations, the zero flag may be set if the result of the operation is zero. Another form of logical data manipulation comes from the rotate instructions such as RR, RL, RRC and RLC which provide a simple means of rotating one bit right or left. Different rotate instructions exist depending on program requirements. Rotate instructions are useful for serial port programming applications where data can be rotated from an internal register into the Carry bit from where it can be examined and the necessary serial bit set high or low. Another application where rotate data operations are used is to implement multiplication and division calculations.

Branches and Control Transfer

Program branching takes the form of either jumps to specified locations using the JMP instruction or to a subroutine using the CALL instruction. They differ in the sense that in the case of a subroutine call, the program must return to the instruction immediately when the subroutine has been carried out. This is done by placing a return instruction RET in the subroutine which will cause the program to jump back to the address right after the CALL instruction. In the case of a JMP instruction, the program simply jumps to the desired location. There is no requirement to jump back to the original jumping off point as in the case of the CALL instruction. One special and extremely useful set of branch instructions are the conditional branches. Here a decision is first made regarding the condition of a certain data memory or individual bits. Depending upon the conditions, the program will continue with the next instruction or skip over it and jump to the following instruction. These instructions are the key to decision making and branching within the program perhaps determined by the condition of certain input switches or by the condition of internal data bits.

Bit Operations

The ability to provide single bit operations on Data Memory is an extremely flexible feature of all Holtek microcontrollers. This feature is especially useful for output port bit programming where individual bits or port pins can be directly set high or low using either the "SET [m].i" or "CLR [m].i" instructions respectively. The feature removes the need for programmers to first read the 8-bit output port, manipulate the input data to ensure that other bits are not changed and then output the port with the correct new data. This read-modify-write process is taken care of automatically when these bit operation instructions are used.

Table Read Operations

Data storage is normally implemented by using registers. However, when working with large amounts of fixed data, the volume involved often makes it inconvenient to store the fixed data in the Data Memory. To overcome this problem, Holtek microcontrollers allow an area of Program Memory to be setup as a table where data can be directly stored. A set of easy to use instructions provides the means by which this fixed data can be referenced and retrieved from the Program Memory.

Other Operations

In addition to the above functional instructions, a range of other instructions also exist such as the "HALT" instruction for Power-down operations and instructions to control the operation of the Watchdog Timer for reliable program operations under extreme electric or electromagnetic environments. For their relevant operations, refer to the functional related sections.

Instruction Set Summary

The following table depicts a summary of the instruction set categorised according to function and can be consulted as a basic instruction reference using the following listed conventions.

Table Conventions

- x: Bits immediate data
- m: Data Memory address
- A: Accumulator
- i: 0~7 number of bits
- addr: Program memory address

Mnemonic	Description	Cycles	Flag Affected
Arithmetic			
ADD A,[m]	Add Data Memory to ACC	1	Z, C, AC, OV
ADDM A,[m]	Add ACC to Data Memory	1 ^{↑Note}	Z, C, AC, OV
ADD A,x	Add immediate data to ACC	1	Z, C, AC, OV
ADC A,[m]	Add Data Memory to ACC with Carry	1	Z, C, AC, OV
ADCM A,[m]	Add ACC to Data memory with Carry	1 ^{↑Note}	Z, C, AC, OV
SUB A,x	Subtract immediate data from the ACC	1	Z, C, AC, OV
SUB A,[m]	Subtract Data Memory from ACC	1	Z, C, AC, OV
SUBM A,[m]	Subtract Data Memory from ACC with result in Data Memory	1 ^{↑Note}	Z, C, AC, OV
SBC A,[m]	Subtract Data Memory from ACC with Carry	1	Z, C, AC, OV
SBCM A,[m]	Subtract Data Memory from ACC with Carry, result in Data Memory	1 ^{↑Note}	Z, C, AC, OV
DAA [m]	Decimal adjust ACC for Addition with result in Data Memory	1 ^{↑Note}	C
Logic Operation			
AND A,[m]	Logical AND Data Memory to ACC	1	Z
OR A,[m]	Logical OR Data Memory to ACC	1	Z
XOR A,[m]	Logical XOR Data Memory to ACC	1	Z
ANDM A,[m]	Logical AND ACC to Data Memory	1 ^{↑Note}	Z
ORM A,[m]	Logical OR ACC to Data Memory	1 ^{↑Note}	Z
XORM A,[m]	Logical XOR ACC to Data Memory	1 ^{↑Note}	Z
AND A,x	Logical AND immediate Data to ACC	1	Z
OR A,x	Logical OR immediate Data to ACC	1	Z
XOR A,x	Logical XOR immediate Data to ACC	1	Z
CPL [m]	Complement Data Memory	1 ^{↑Note}	Z
CPLA [m]	Complement Data Memory with result in ACC	1	Z
Increment & Decrement			
INCA [m]	Increment Data Memory with result in ACC	1	Z
INC [m]	Increment Data Memory	1 ^{↑Note}	Z
DECA [m]	Decrement Data Memory with result in ACC	1	Z
DEC [m]	Decrement Data Memory	1 ^{↑Note}	Z
Rotate			
RRA [m]	Rotate Data Memory right with result in ACC	1	None
RR [m]	Rotate Data Memory right	1 ^{↑Note}	None
RRCA [m]	Rotate Data Memory right through Carry with result in ACC	1	C
RRC [m]	Rotate Data Memory right through Carry	1 ^{↑Note}	C
RLA [m]	Rotate Data Memory left with result in ACC	1	None
RL [m]	Rotate Data Memory left	1 ^{↑Note}	None
RLCA [m]	Rotate Data Memory left through Carry with result in ACC	1	C
RLC [m]	Rotate Data Memory left through Carry	1 ^{↑Note}	C

Mnemonic	Description	Cycles	Flag Affected
Data Move			
MOV A,[m]	Move Data Memory to ACC	1	None
MOV [m],A	Move ACC to Data Memory	1 ^{Note}	None
MOV A,x	Move immediate data to ACC	1	None
Bit Operation			
CLR [m].i	Clear bit of Data Memory	1 ^{Note}	None
SET [m].i	Set bit of Data Memory	1 ^{Note}	None
Branch			
JMP addr	Jump unconditionally	2	None
SZ [m]	Skip if Data Memory is zero	1 ^{Note}	None
SZA [m]	Skip if Data Memory is zero with data movement to ACC	1 ^{Note}	None
SZ [m].i	Skip if bit i of Data Memory is zero	1 ^{Note}	None
SNZ [m].i	Skip if bit i of Data Memory is not zero	1 ^{Note}	None
SIZ [m]	Skip if increment Data Memory is zero	1 ^{Note}	None
SDZ [m]	Skip if decrement Data Memory is zero	1 ^{Note}	None
SIZA [m]	Skip if increment Data Memory is zero with result in ACC	1 ^{Note}	None
SDZA [m]	Skip if decrement Data Memory is zero with result in ACC	1 ^{Note}	None
CALL addr	Subroutine call	2	None
RET	Return from subroutine	2	None
RET A,x	Return from subroutine and load immediate data to ACC	2	None
RETI	Return from interrupt	2	None
Table Read			
TABRD [m]	Read table to TBLH and Data Memory	2 ^{Note}	None
TABRDL [m]	Read table (last page) to TBLH and Data Memory	2 ^{Note}	None
Miscellaneous			
NOP	No operation	1	None
CLR [m]	Clear Data Memory	1 ^{Note}	None
SET [m]	Set Data Memory	1 ^{Note}	None
CLR WDT	Clear Watchdog Timer	1	TO, PDF
CLR WDT1	Pre-clear Watchdog Timer	1	TO, PDF
CLR WDT2	Pre-clear Watchdog Timer	1	TO, PDF
SWAP [m]	Swap nibbles of Data Memory	1 ^{Note}	None
SWAPA [m]	Swap nibbles of Data Memory with result in ACC	1	None
HALT	Enter power down mode	1	TO, PDF

Note: 1. For skip instructions, if the result of the comparison involves a skip then two cycles are required, if no skip takes place only one cycle is required.

2. Any instruction which changes the contents of the PCL will also require 2 cycles for execution.

3. For the "CLR WDT1" and "CLR WDT2" instructions the TO and PDF flags may be affected by the execution status. The TO and PDF flags are cleared after both "CLR WDT1" and "CLR WDT2" instructions are consecutively executed. Otherwise the TO and PDF flags remain unchanged.

Instruction Definition

ADC A,[m]	Add Data Memory to ACC with Carry
Description	The contents of the specified Data Memory, Accumulator and the carry flag are added. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC + [m] + C$
Affected flag(s)	OV, Z, AC, C
ADCM A,[m]	Add ACC to Data Memory with Carry
Description	The contents of the specified Data Memory, Accumulator and the carry flag are added. The result is stored in the specified Data Memory.
Operation	$[m] \leftarrow ACC + [m] + C$
Affected flag(s)	OV, Z, AC, C
ADD A,[m]	Add Data Memory to ACC
Description	The contents of the specified Data Memory and the Accumulator are added. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC + [m]$
Affected flag(s)	OV, Z, AC, C
ADD A,x	Add immediate data to ACC
Description	The contents of the Accumulator and the specified immediate data are added. The result is stored in the Accumulator.
Operation	$AC \leftarrow ACC + x$
Affected flag(s)	OV, Z, AC, C
ADDM A,[m]	Add ACC to Data Memory
Description	The contents of the specified Data Memory and the Accumulator are added. The result is stored in the specified Data Memory.
Operation	$[m] \leftarrow ACC + [m]$
Affected flag(s)	OV, Z, AC, C
AND A,[m]	Logical AND Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical AND operation. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC \text{ "AND" } [m]$
Affected flag(s)	Z
AND A,x	Logical AND immediate data to ACC
Description	Data in the Accumulator and the specified immediate data perform a bitwise logical AND operation. The result is stored in the Accumulator.
Operation	$ACC \leftarrow ACC \text{ "AND" } x$
Affected flag(s)	Z
ANDM A,[m]	Logical AND ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical AND operation. The result is stored in the Data Memory.
Operation	$[m] \leftarrow ACC \text{ "AND" } [m]$
Affected flag(s)	Z

CALL addr	Subroutine call
Description	Unconditionally calls a subroutine at the specified address. The Program Counter then increments by 1 to obtain the address of the next instruction which is then pushed onto the stack. The specified address is then loaded and the program continues execution from this new address. As this instruction requires an additional operation, it is a two cycle instruction.
Operation	Stack \leftarrow Program Counter + 1 Program Counter \leftarrow addr
Affected flag(s)	None
CLR [m]	Clear Data Memory
Description	Each bit of the specified Data Memory is cleared to 0.
Operation	[m] \leftarrow 00H
Affected flag(s)	None
CLR [m].i	Clear bit of Data Memory
Description	Bit i of the specified Data Memory is cleared to 0.
Operation	[m].i \leftarrow 0
Affected flag(s)	None
CLR WDT	Clear Watchdog Timer
Description	The TO, PDF flags and the WDT are all cleared.
Operation	WDT cleared TO \leftarrow 0 PDF \leftarrow 0
Affected flag(s)	TO, PDF
CLR WDT1	Pre-clear Watchdog Timer
Description	The TO, PDF flags and the WDT are all cleared. Note that this instruction works in conjunction with CLR WDT2 and must be executed alternately with CLR WDT2 to have effect. Repetitively executing this instruction without alternately executing CLR WDT2 will have no effect.
Operation	WDT cleared TO \leftarrow 0 PDF \leftarrow 0
Affected flag(s)	TO, PDF
CLR WDT2	Pre-clear Watchdog Timer
Description	The TO, PDF flags and the WDT are all cleared. Note that this instruction works in conjunction with CLR WDT1 and must be executed alternately with CLR WDT1 to have effect. Repetitively executing this instruction without alternately executing CLR WDT1 will have no effect.
Operation	WDT cleared TO \leftarrow 0 PDF \leftarrow 0
Affected flag(s)	TO, PDF

CPL [m]	Complement Data Memory
Description	Each bit of the specified Data Memory is logically complemented (1's complement). Bits which previously contained a 1 are changed to 0 and vice versa.
Operation	$[m] \leftarrow \overline{[m]}$
Affected flag(s)	Z
CPLA [m]	Complement Data Memory with result in ACC
Description	Each bit of the specified Data Memory is logically complemented (1's complement). Bits which previously contained a 1 are changed to 0 and vice versa. The complemented result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC \leftarrow \overline{[m]}$
Affected flag(s)	Z
DAA [m]	Decimal-Adjust ACC for addition with result in Data Memory
Description	Convert the contents of the Accumulator value to a BCD (Binary Coded Decimal) value resulting from the previous addition of two BCD variables. If the low nibble is greater than 9 or if AC flag is set, then a value of 6 will be added to the low nibble. Otherwise the low nibble remains unchanged. If the high nibble is greater than 9 or if the C flag is set, then a value of 6 will be added to the high nibble. Essentially, the decimal conversion is performed by adding 00H, 06H, 60H or 66H depending on the Accumulator and flag conditions. Only the C flag may be affected by this instruction which indicates that if the original BCD sum is greater than 100, it allows multiple precision decimal addition.
Operation	$[m] \leftarrow ACC + 00H$ or $[m] \leftarrow ACC + 06H$ or $[m] \leftarrow ACC + 60H$ or $[m] \leftarrow ACC + 66H$
Affected flag(s)	C
DEC [m]	Decrement Data Memory
Description	Data in the specified Data Memory is decremented by 1.
Operation	$[m] \leftarrow [m] - 1$
Affected flag(s)	Z
DECA [m]	Decrement Data Memory with result in ACC
Description	Data in the specified Data Memory is decremented by 1. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	$ACC \leftarrow [m] - 1$
Affected flag(s)	Z

HALT	Enter power down mode
Description	This instruction stops the program execution and turns off the system clock. The contents of the Data Memory and registers are retained. The WDT and prescaler are cleared. The power down flag PDF is set and the WDT time-out flag TO is cleared.
Operation	TO ← 0 PDF ← 0
Affected flag(s)	TO, PDF
INC [m]	Increment Data Memory
Description	Data in the specified Data Memory is incremented by 1.
Operation	[m] ← [m]+1
Affected flag(s)	Z
INCA [m]	Increment Data Memory with result in ACC
Description	Data in the specified Data Memory is incremented by 1. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	ACC ← [m]+1
Affected flag(s)	Z
JMP addr	Jump unconditionally
Description	The contents of the Program Counter are replaced with the specified address. Program execution then continues from this new address. As this requires the insertion of a dummy instruction while the new address is loaded, it is a two cycle instruction.
Operation	Program Counter ← addr
Affected flag(s)	None
MOV A,[m]	Move Data Memory to ACC
Description	The contents of the specified Data Memory are copied to the Accumulator.
Operation	ACC ← [m]
Affected flag(s)	None
MOV A,x	Move immediate data to ACC
Description	The immediate data specified is loaded into the Accumulator.
Operation	ACC ← x
Affected flag(s)	None
MOV [m],A	Move ACC to Data Memory
Description	The contents of the Accumulator are copied to the specified Data Memory.
Operation	[m] ← ACC
Affected flag(s)	None
NOP	No operation
Description	No operation is performed. Execution continues with the next instruction.
Operation	No operation
Affected flag(s)	None

OR A,[m]	Logical OR Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical OR operation. The result is stored in the Accumulator.
Operation	ACC ← ACC "OR" [m]
Affected flag(s)	Z
OR A,x	Logical OR immediate data to ACC
Description	Data in the Accumulator and the specified immediate data perform a bitwise logical OR operation. The result is stored in the Accumulator.
Operation	ACC ← ACC "OR" x
Affected flag(s)	Z
ORM A,[m]	Logical OR ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical OR operation. The result is stored in the Data Memory.
Operation	[m] ← ACC "OR" [m]
Affected flag(s)	Z
RET	Return from subroutine
Description	The Program Counter is restored from the stack. Program execution continues at the restored address.
Operation	Program Counter ← Stack
Affected flag(s)	None
RET A,x	Return from subroutine and load immediate data to ACC
Description	The Program Counter is restored from the stack and the Accumulator loaded with the specified immediate data. Program execution continues at the restored address.
Operation	Program Counter ← Stack ACC ← x
Affected flag(s)	None
RETI	Return from interrupt
Description	The Program Counter is restored from the stack and the interrupts are re-enabled by setting the EMI bit. EMI is the master interrupt global enable bit. If an interrupt was pending when the RETI instruction is executed, the pending Interrupt routine will be processed before returning to the main program.
Operation	Program Counter ← Stack EMI ← 1
Affected flag(s)	None
RL [m]	Rotate Data Memory left
Description	The contents of the specified Data Memory are rotated left by 1 bit with bit 7 rotated into bit 0.
Operation	[m].(i+1) ← [m].i; (i = 0~6) [m].0 ← [m].7
Affected flag(s)	None

RLA [m]	Rotate Data Memory left with result in ACC
Description	The contents of the specified Data Memory are rotated left by 1 bit with bit 7 rotated into bit 0. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	ACC.(i+1) ← [m].i; (i = 0~6) ACC.0 ← [m].7
Affected flag(s)	None
RLC [m]	Rotate Data Memory left through Carry
Description	The contents of the specified Data Memory and the carry flag are rotated left by 1 bit. Bit 7 replaces the Carry bit and the original carry flag is rotated into bit 0.
Operation	[m].(i+1) ← [m].i; (i = 0~6) [m].0 ← C C ← [m].7
Affected flag(s)	C
RLCA [m]	Rotate Data Memory left through Carry with result in ACC
Description	Data in the specified Data Memory and the carry flag are rotated left by 1 bit. Bit 7 replaces the Carry bit and the original carry flag is rotated into the bit 0. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	ACC.(i+1) ← [m].i; (i = 0~6) ACC.0 ← C C ← [m].7
Affected flag(s)	C
RR [m]	Rotate Data Memory right
Description	The contents of the specified Data Memory are rotated right by 1 bit with bit 0 rotated into bit 7.
Operation	[m].i ← [m].(i+1); (i = 0~6) [m].7 ← [m].0
Affected flag(s)	None
RRA [m]	Rotate Data Memory right with result in ACC
Description	Data in the specified Data Memory and the carry flag are rotated right by 1 bit with bit 0 rotated into bit 7. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	ACC.i ← [m].(i+1); (i = 0~6) ACC.7 ← [m].0
Affected flag(s)	None
RRC [m]	Rotate Data Memory right through Carry
Description	The contents of the specified Data Memory and the carry flag are rotated right by 1 bit. Bit 0 replaces the Carry bit and the original carry flag is rotated into bit 7.
Operation	[m].i ← [m].(i+1); (i = 0~6) [m].7 ← C C ← [m].0
Affected flag(s)	C

RRCA [m]	Rotate Data Memory right through Carry with result in ACC
Description	Data in the specified Data Memory and the carry flag are rotated right by 1 bit. Bit 0 replaces the Carry bit and the original carry flag is rotated into bit 7. The rotated result is stored in the Accumulator and the contents of the Data Memory remain unchanged.
Operation	$ACC.i \leftarrow [m].(i+1); (i = 0\sim 6)$ $ACC.7 \leftarrow C$ $C \leftarrow [m].0$
Affected flag(s)	C
SBC A,[m]	Subtract Data Memory from ACC with Carry
Description	The contents of the specified Data Memory and the complement of the carry flag are subtracted from the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - [m] - \bar{C}$
Affected flag(s)	OV, Z, AC, C
SBCM A,[m]	Subtract Data Memory from ACC with Carry and result in Data Memory
Description	The contents of the specified Data Memory and the complement of the carry flag are subtracted from the Accumulator. The result is stored in the Data Memory. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - [m] - \bar{C}$
Affected flag(s)	OV, Z, AC, C
SDZ [m]	Skip if decrement Data Memory is 0
Description	The contents of the specified Data Memory are first decremented by 1. If the result is 0 the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$[m] \leftarrow [m] - 1$ Skip if $[m] = 0$
Affected flag(s)	None
SDZA [m]	Skip if decrement Data Memory is zero with result in ACC
Description	The contents of the specified Data Memory are first decremented by 1. If the result is 0, the following instruction is skipped. The result is stored in the Accumulator but the specified Data Memory contents remain unchanged. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0, the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m] - 1$
Affected flag(s)	Skip if $ACC = 0$ None

SET [m]	Set Data Memory
Description	Each bit of the specified Data Memory is set to 1.
Operation	$[m] \leftarrow FFH$
Affected flag(s)	None
SET [m].i	Set bit of Data Memory
Description	Bit i of the specified Data Memory is set to 1.
Operation	$[m].i \leftarrow 1$
Affected flag(s)	None
SIZ [m]	Skip if increment Data Memory is 0
Description	The contents of the specified Data Memory are first incremented by 1. If the result is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$[m] \leftarrow [m] + 1$ Skip if $[m] = 0$
Affected flag(s)	None
SIZA [m]	Skip if increment Data Memory is zero with result in ACC
Description	The contents of the specified Data Memory are first incremented by 1. If the result is 0, the following instruction is skipped. The result is stored in the Accumulator but the specified Data Memory contents remain unchanged. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m] + 1$ Skip if $ACC = 0$
Affected flag(s)	None
SNZ [m].i	Skip if bit i of Data Memory is not 0
Description	If bit i of the specified Data Memory is not 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is 0 the program proceeds with the following instruction.
Operation	Skip if $[m].i \neq 0$
Affected flag(s)	None
SUB A,[m]	Subtract Data Memory from ACC
Description	The specified Data Memory is subtracted from the contents of the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - [m]$
Affected flag(s)	OV, Z, AC, C

SUBM A,[m]	Subtract Data Memory from ACC with result in Data Memory
Description	The specified Data Memory is subtracted from the contents of the Accumulator. The result is stored in the Data Memory. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$[m] \leftarrow ACC - [m]$
Affected flag(s)	OV, Z, AC, C
SUB A,x	Subtract immediate data from ACC
Description	The immediate data specified by the code is subtracted from the contents of the Accumulator. The result is stored in the Accumulator. Note that if the result of subtraction is negative, the C flag will be cleared to 0, otherwise if the result is positive or zero, the C flag will be set to 1.
Operation	$ACC \leftarrow ACC - x$
Affected flag(s)	OV, Z, AC, C
SWAP [m]	Swap nibbles of Data Memory
Description	The low-order and high-order nibbles of the specified Data Memory are interchanged.
Operation	$[m].3 \sim [m].0 \leftrightarrow [m].7 \sim [m].4$
Affected flag(s)	None
SWAPA [m]	Swap nibbles of Data Memory with result in ACC
Description	The low-order and high-order nibbles of the specified Data Memory are interchanged. The result is stored in the Accumulator. The contents of the Data Memory remain unchanged.
Operation	$ACC.3 \sim ACC.0 \leftarrow [m].7 \sim [m].4$ $ACC.7 \sim ACC.4 \leftarrow [m].3 \sim [m].0$
Affected flag(s)	None
SZ [m]	Skip if Data Memory is 0
Description	If the contents of the specified Data Memory is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	Skip if $[m] = 0$
Affected flag(s)	None
SZA [m]	Skip if Data Memory is 0 with data movement to ACC
Description	The contents of the specified Data Memory are copied to the Accumulator. If the value is zero, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0 the program proceeds with the following instruction.
Operation	$ACC \leftarrow [m]$ Skip if $[m] = 0$
Affected flag(s)	None

SZ [m].i	Skip if bit i of Data Memory is 0
Description	If bit i of the specified Data Memory is 0, the following instruction is skipped. As this requires the insertion of a dummy instruction while the next instruction is fetched, it is a two cycle instruction. If the result is not 0, the program proceeds with the following instruction.
Operation	Skip if [m].i = 0
Affected flag(s)	None
TABRD [m]	Read table to TBLH and Data Memory
Description	The program code addressed by the table pointer (TBHP and TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
TABRDL [m]	Read table (last page) to TBLH and Data Memory
Description	The low byte of the program code (last page) addressed by the table pointer (TBLP) is moved to the specified Data Memory and the high byte moved to TBLH.
Operation	[m] ← program code (low byte) TBLH ← program code (high byte)
Affected flag(s)	None
XOR A,[m]	Logical XOR Data Memory to ACC
Description	Data in the Accumulator and the specified Data Memory perform a bitwise logical XOR operation. The result is stored in the Accumulator.
Operation	ACC ← ACC "XOR" [m]
Affected flag(s)	Z
XORM A,[m]	Logical XOR ACC to Data Memory
Description	Data in the specified Data Memory and the Accumulator perform a bitwise logical XOR operation. The result is stored in the Data Memory.
Operation	[m] ← ACC "XOR" [m]
Affected flag(s)	Z
XOR A,x	Logical XOR immediate data to ACC
Description	Data in the Accumulator and the specified immediate data perform a bitwise logical XOR operation. The result is stored in the Accumulator.
Operation	ACC ← ACC "XOR" x
Affected flag(s)	Z

Package Information

Note that the package information provided here is for consultation purposes only. As this information may be updated at regular intervals users are reminded to consult the Holtek website (<http://www.holtek.com.tw/english/literature/package.pdf>) for the latest version of the package information.

16-pin DIP (300mil) Outline Dimensions

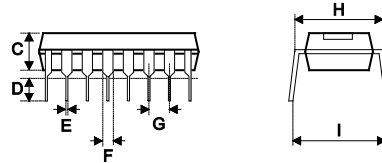
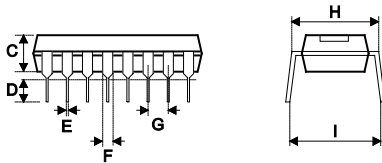
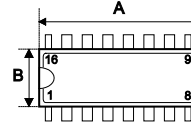
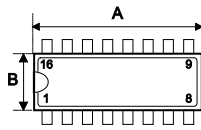


Fig1. Full Lead Packages

Fig2. 1/2 Lead Packages

MS-001d (see fig1)

Symbol	Dimensions in inch		
	Min.	Nom.	Max.
A	0.780	—	0.880
B	0.240	—	0.280
C	0.115	—	0.195
D	0.115	—	0.150
E	0.014	—	0.022
F	0.045	—	0.070
G	—	0.100	—
H	0.300	—	0.325
I	—	0.430	—

Symbol	Dimensions in mm		
	Min.	Nom.	Max.
A	19.81	—	22.35
B	6.10	—	7.11
C	2.92	—	4.95
D	2.92	—	3.81
E	0.36	—	0.56
F	1.14	—	1.78
G	—	2.54	—
H	7.62	—	8.26
I	—	10.92	—

MS-001d (see fig2)

Symbol	Dimensions in inch		
	Min.	Nom.	Max.
A	0.735	—	0.775
B	0.240	—	0.280
C	0.115	—	0.195
D	0.115	—	0.150
E	0.014	—	0.022
F	0.045	—	0.070
G	—	0.100	—
H	0.300	—	0.325
I	—	0.430	—

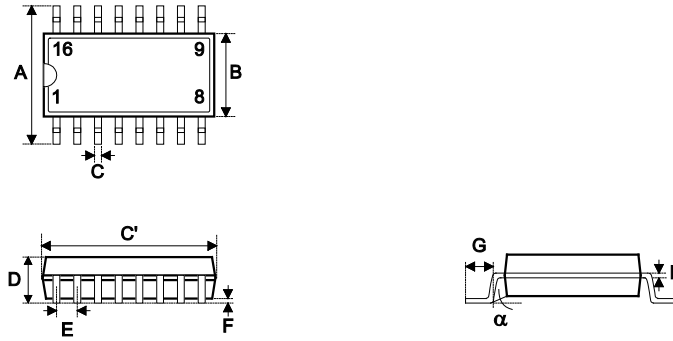
Symbol	Dimensions in mm		
	Min.	Nom.	Max.
A	18.67	—	19.69
B	6.10	—	7.11
C	2.92	—	4.95
D	2.92	—	3.81
E	0.36	—	0.56
F	1.14	—	1.78
G	—	2.54	—
H	7.62	—	8.26
I	—	10.92	—

MO-095a (see fig2)

Symbol	Dimensions in inch		
	Min.	Nom.	Max.
A	0.745	—	0.785
B	0.275	—	0.295
C	0.120	—	0.150
D	0.110	—	0.150
E	0.014	—	0.022
F	0.045	—	0.060
G	—	0.100	—
H	0.300	—	0.325
I	—	0.430	—

Symbol	Dimensions in mm		
	Min.	Nom.	Max.
A	18.92	—	19.94
B	6.99	—	7.49
C	3.05	—	3.81
D	2.79	—	3.81
E	0.36	—	0.56
F	1.14	—	1.52
G	—	2.54	—
H	7.62	—	8.26
I	—	10.92	—

16-pin NSOP (150mil) Outline Dimensions

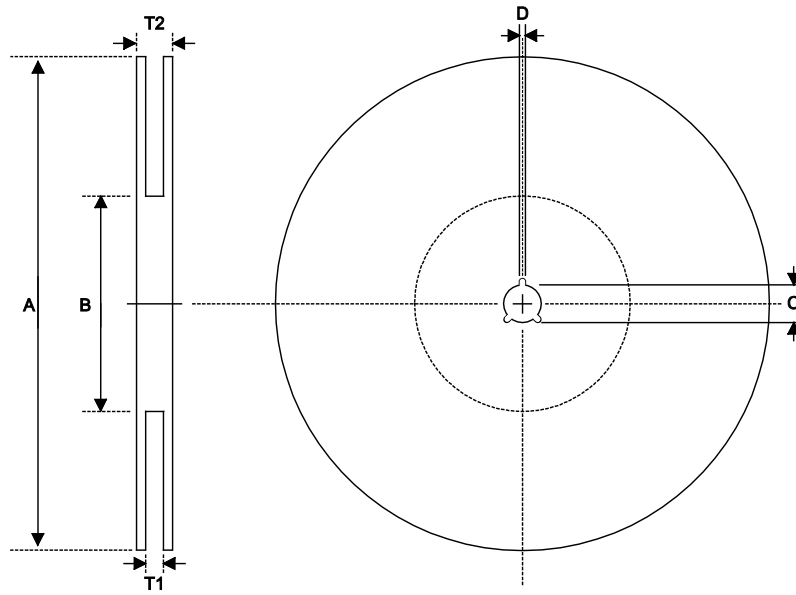


MS-012

Symbol	Dimensions in inch		
	Min.	Nom.	Max.
A	0.228	—	0.244
B	0.150	—	0.157
C	0.012	—	0.020
C'	0.386	—	0.402
D	—	—	0.069
E	—	0.050	—
F	0.004	—	0.010
G	0.016	—	0.050
H	0.007	—	0.010
α	0°	—	8°

Symbol	Dimensions in mm		
	Min.	Nom.	Max.
A	5.79	—	6.20
B	3.81	—	3.99
C	0.30	—	0.51
C'	9.80	—	10.21
D	—	—	1.75
E	—	1.27	—
F	0.10	—	0.25
G	0.41	—	1.27
H	0.18	—	0.25
α	0°	—	8°

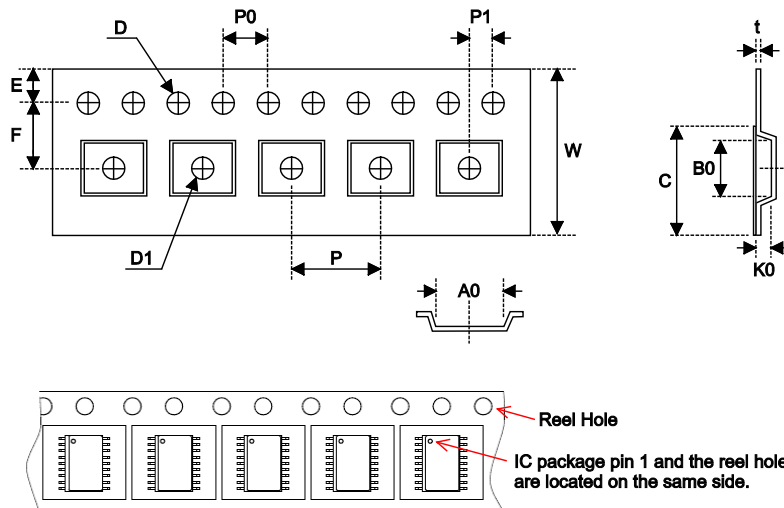
Reel Dimensions



SOP 16NSOP (150mil)

Symbol	Description	Dimensions in mm
A	Reel Outer Diameter	330.0±1.0
B	Reel Inner Diameter	100.0±1.5
C	Spindle Hole Diameter	13.0 ^{+0.5/-0.2}
D	Key Slit Width	2.0±0.5
T1	Space Between Flange	16.8 ^{+0.3/-0.2}
T2	Reel Thickness	22.2±0.2

Carrier Tape Dimensions



16-pin NSOP (150mil)

Symbol	Description	Symbol Description Dimensions in mm
W	Carrier Tape Width	16.0±0.3
P	Cavity Pitch	8.0±0.1
E	Perforation Position	1.75±0.1
F	Cavity to Perforation (Width Direction)	7.5±0.1
D	Perforation Diameter	1.55 ^{+0.10/-0.00}
D1	Cavity Hole Diameter	1.50 ^{+0.25/-0.00}
P0	Perforation Pitch	4.0±0.1
P1	Cavity to Perforation (Length Direction)	2.0±0.1
A0	Cavity Length	6.5±0.1
B0	Cavity Width	10.3±0.1
K0	Cavity Depth	2.1±0.1
t	Carrier Tape Thickness	0.30±0.05
C	Cover Tape Width	13.3±0.1

Holtek Semiconductor Inc. (Headquarters)

No.3, Creation Rd. II, Science Park, Hsinchu, Taiwan
Tel: 886-3-563-1999
Fax: 886-3-563-1189
<http://www.holtek.com.tw>

Holtek Semiconductor Inc. (Taipei Sales Office)

4F-2, No. 3-2, YuanQu St., Nankang Software Park, Taipei 115, Taiwan
Tel: 886-2-2655-7070
Fax: 886-2-2655-7373
Fax: 886-2-2655-7383 (International sales hotline)

Holtek Semiconductor (China) Inc. (Dongguan Sales Office)

Building No.10, Xinzhu Court, (No.1 Headquarters), 4 Cuizhu Road, Songshan Lake, Dongguan, China 523808
Tel: 86-769-2626-1300
Fax: 86-769-2626-1311, 86-769-2626-1322

Holtek Semiconductor (USA), Inc. (North America Sales Office)

46729 Fremont Blvd., Fremont, CA 94538, USA
Tel: 1-510-252-9880
Fax: 1-510-252-9885
<http://www.holtek.com>

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